

Research Challenges in Orchestration Synthesis

Davide Basile, Maurice H. ter Beek

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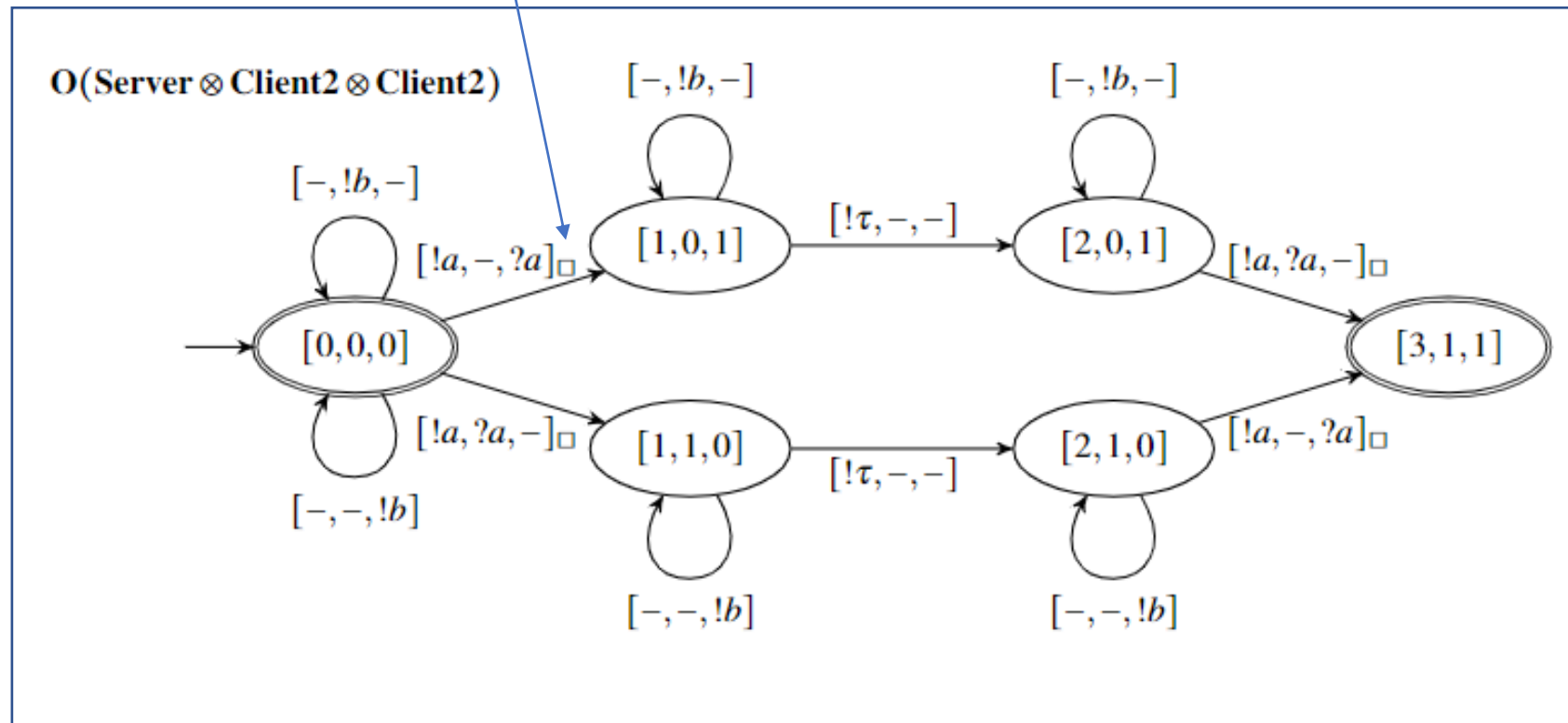
Overview

- Introduction to contract automata
- Orchestration synthesis algorithm
- Research Challenges
- Conclusion

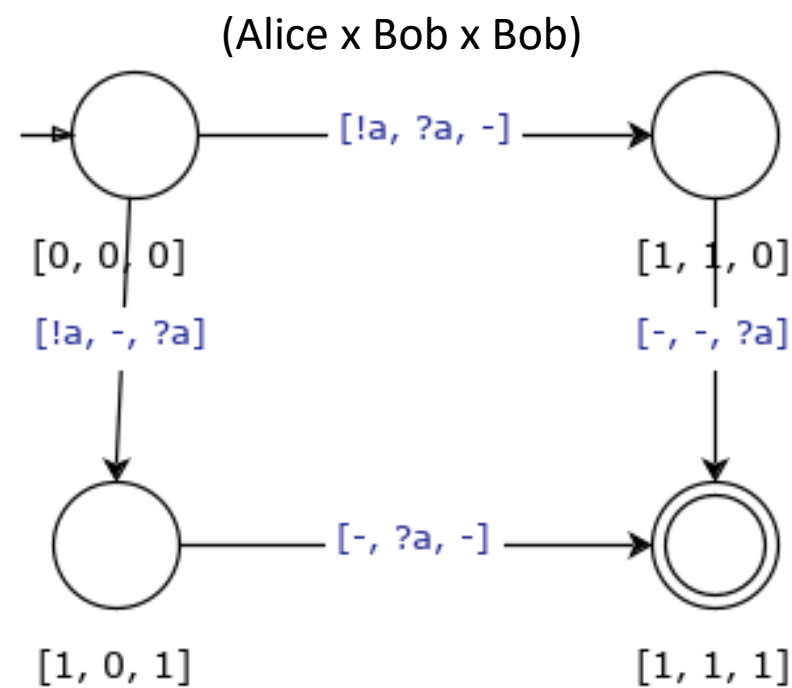
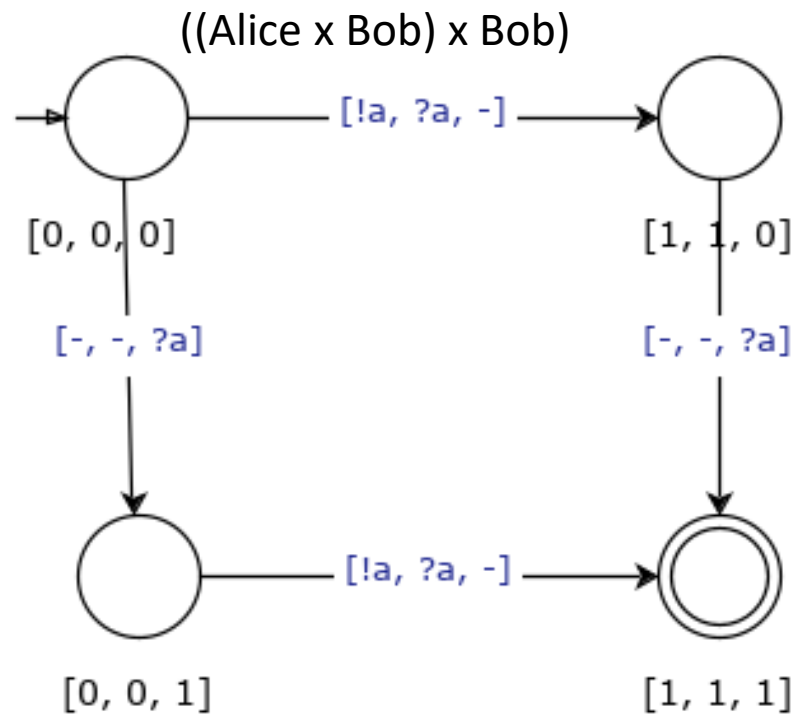
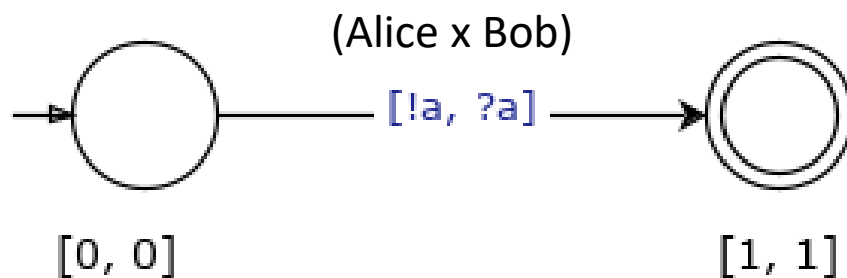
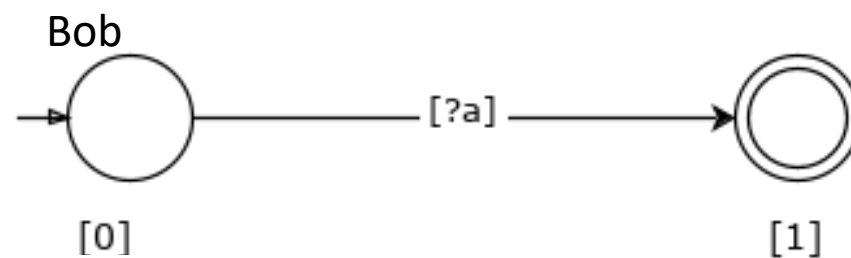
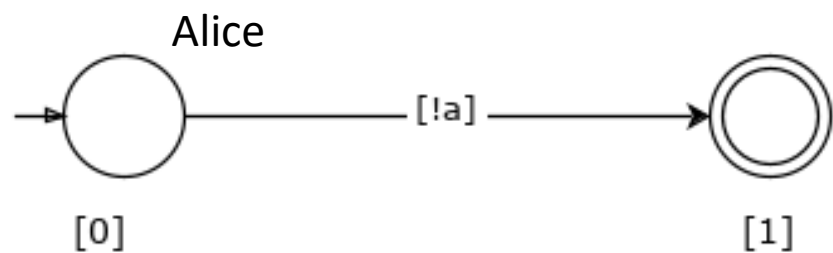
Contract Automata

- Behavioural contracts: used to specify the interactions of services
- Contract automata: a dialect of finite-state automata modelling behavioural contracts
- Introduced almost 10 years ago
 - Basile, D., Degano, P. and Ferrari, G.L. Automata for analysing service contracts. In *TGC 2014*.
 - Basile, D., Degano, P., Ferrari, G.L. and Tuosto, E. From orchestration to choreography through contract automata. In *ICE 2014*.
- We will present challenges in the orchestration synthesis of contract automata

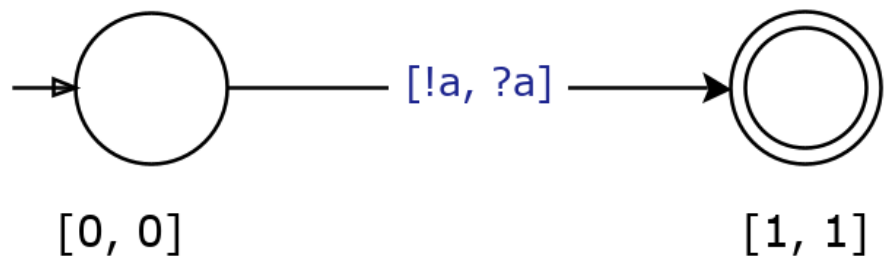
- Contract automata are FSA enhanced with:
 - Partitioned alphabet of actions:
 - offers $!a$ (A^o) and requests $?a$ (A^r)
 - special idle action $-$ (not in $A^o \cup A^r$)
 - rank : the number of services in the contract
 - Transitions partitioned into *optional* and *necessary*
 - States are list of basic states
 - Labels are list of actions and are constrained to be:
 - offers: $(-, -, -, !a)$
 - requests: $(-, ?a, -, -)$
 - matches: $(-, ?a, -, !a)$
(only between two)
 - $\text{size}(\text{list}) = \text{rank}$



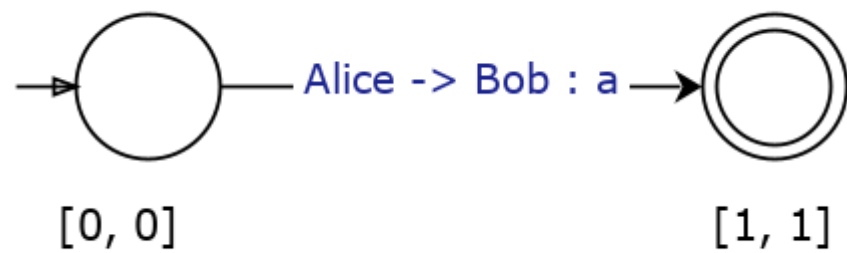
Compositionality, "Partiality"



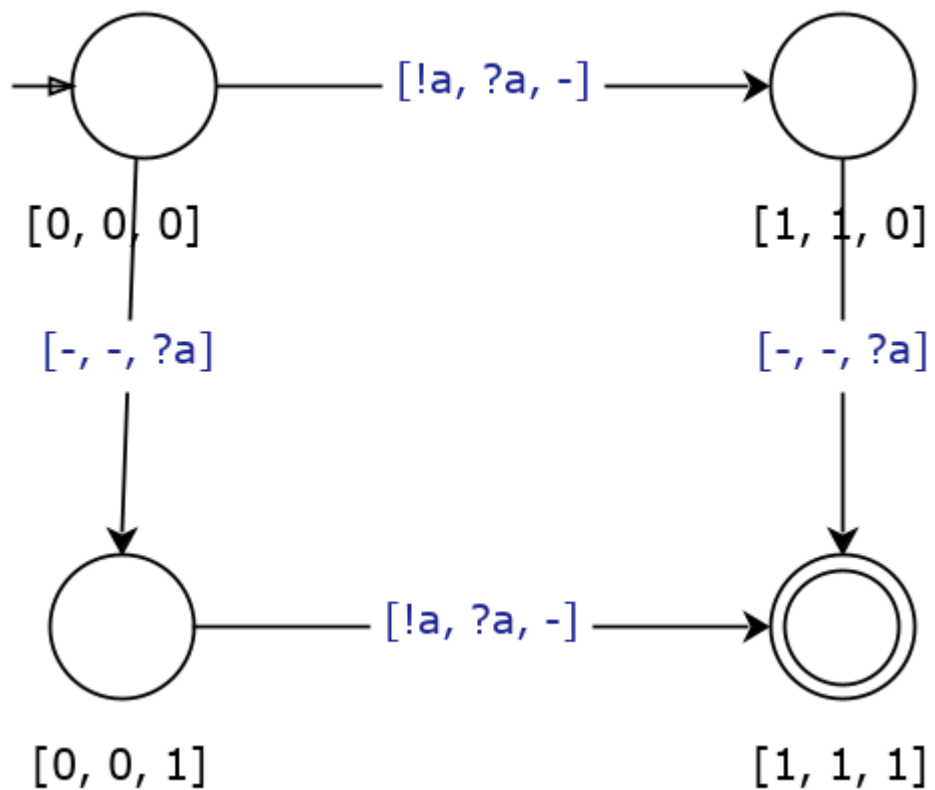
(Alice x Bob)



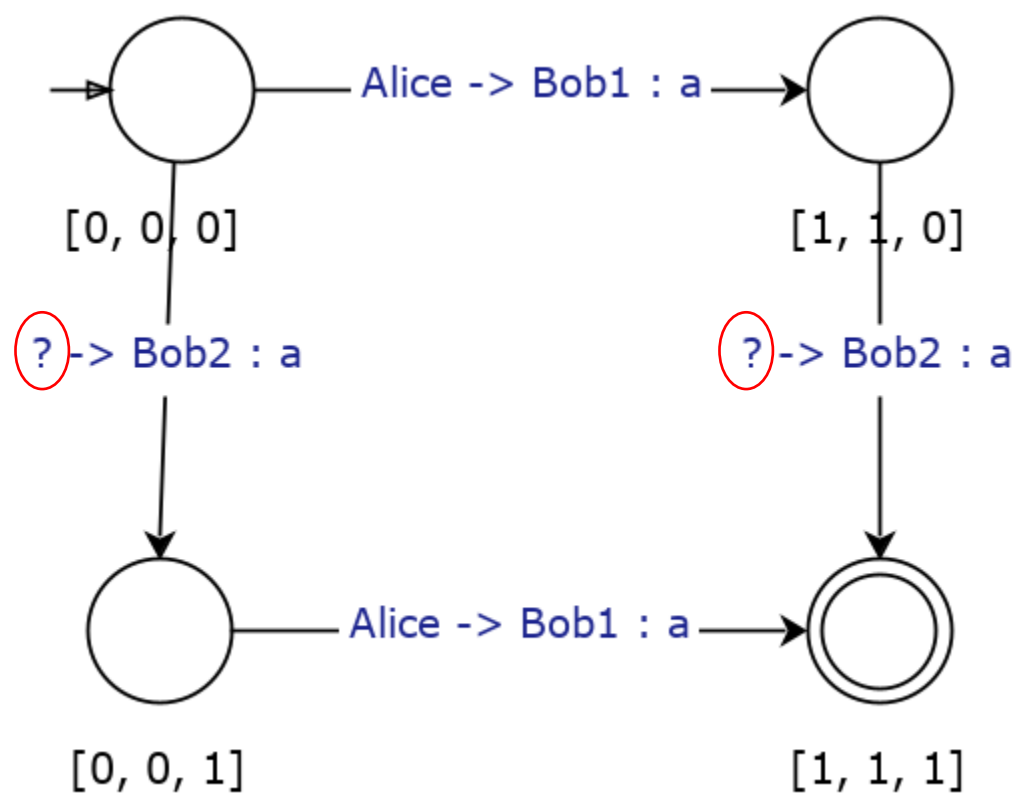
=



((Alice x Bob₁) x Bob₂)



=



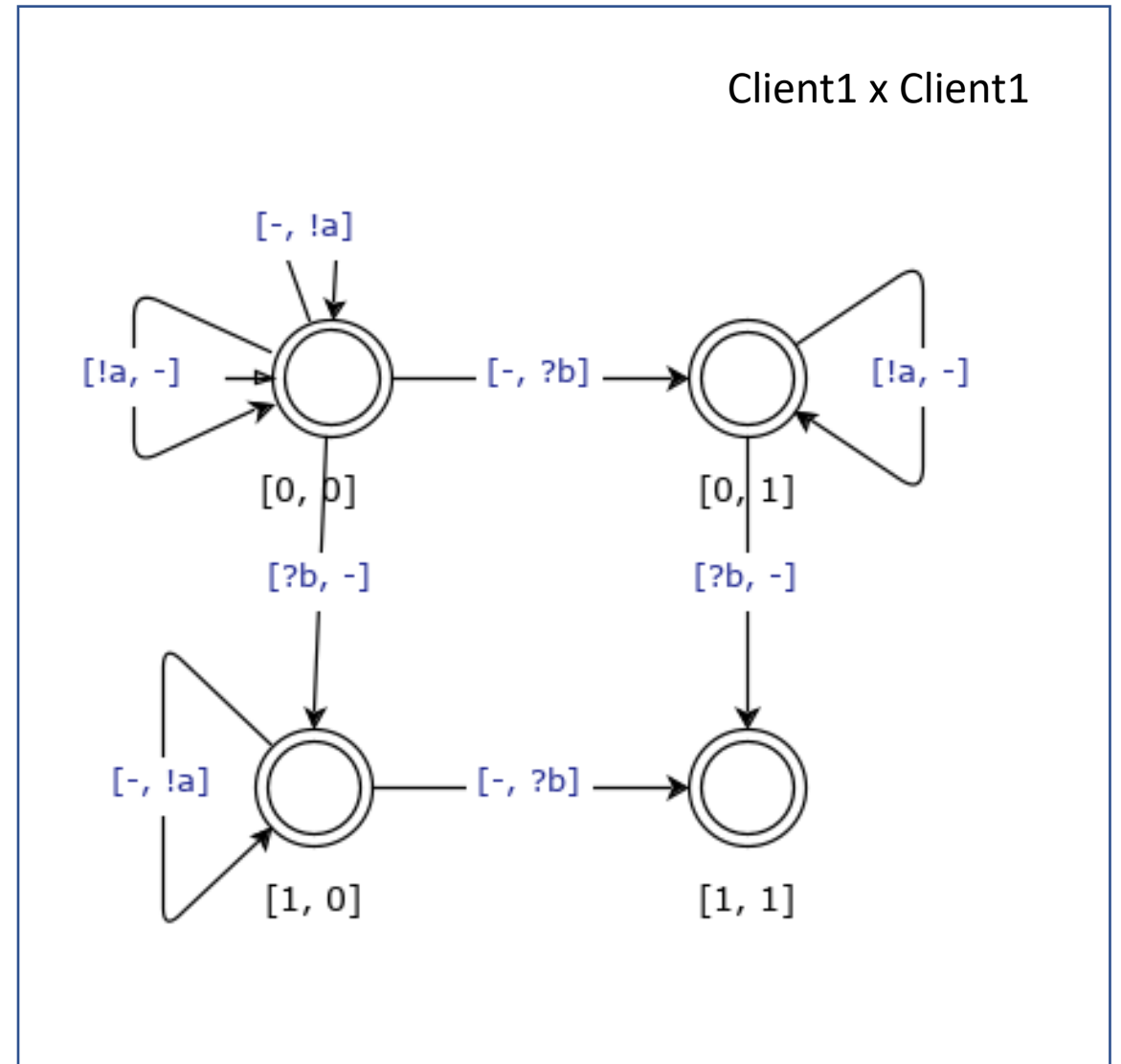
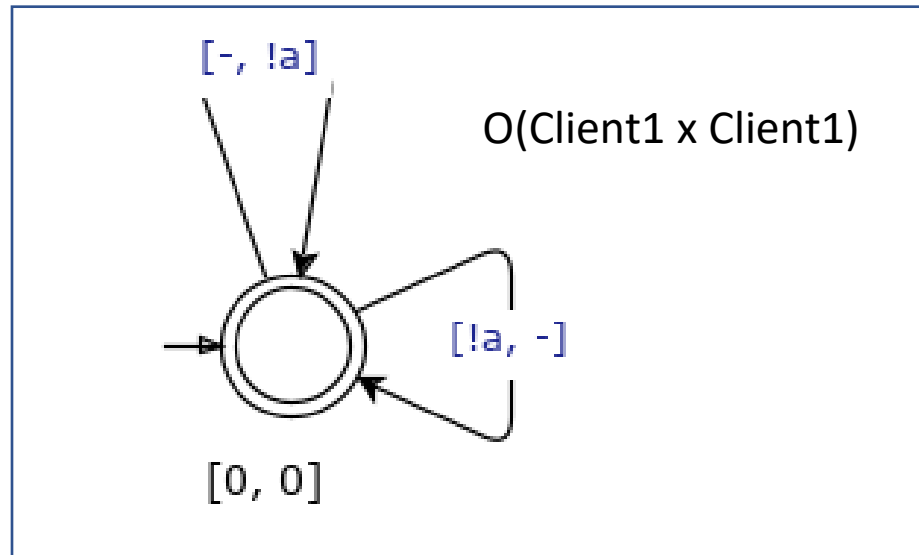
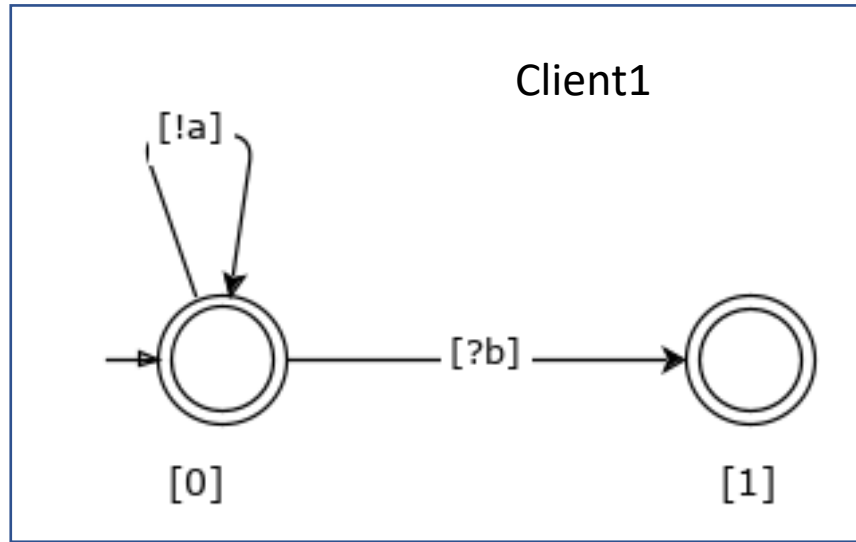
Orchestration Synthesis

- The orchestration is computed using the synthesis of the most permissive controller (mpc) from Supervisory Control Theory (SCT) for discrete event systems
- The orchestrator adapts already existing services by blocking some of their actions to enforce *agreement* among the parties
 - $(!a,?a,-)(!a,-,-)$ trace is in agreement (all requests are matched)
 - $(!a,?a,-)(-,?a,-)$ trace not in agreement ($(-,?a,-)$ is a request label)
- The orchestrator is *abstracted* away
- The synthesised orchestration assumes the presence of an orchestrator invoking the actions of the services

Properties of the orchestration

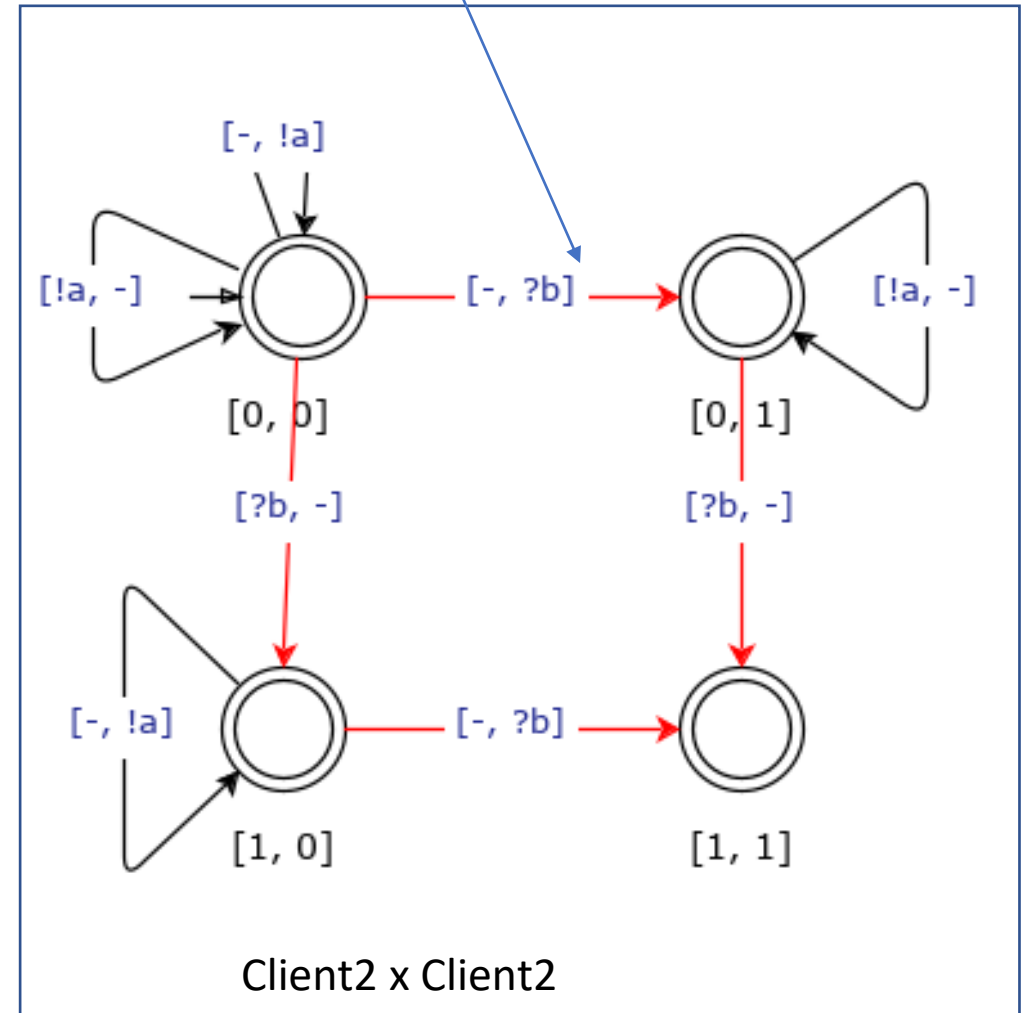
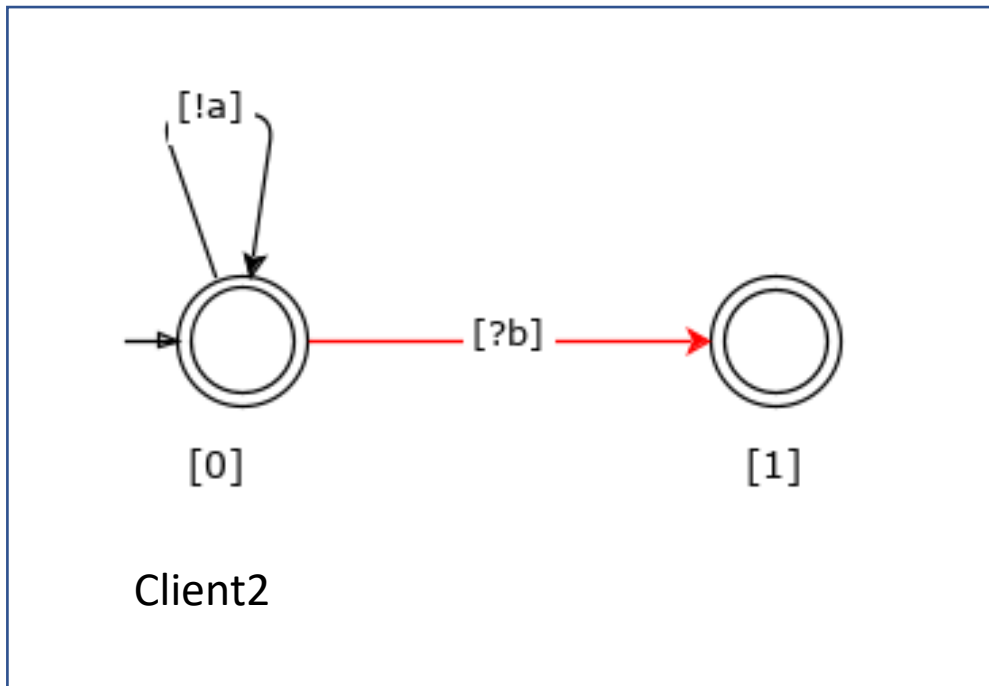
- The synthesis computes a refinement of the composition
 - Minimal fixed point algorithm
- SCT distinguishes between controllable and uncontrollable actions
- Only controllable actions can be removed (*controllability*)
- Final states (a.k.a. marked states) are always reachable (*non-blocking*)
- Forbidden states are never traversed (*safety*)
 - States with outgoing uncontrollable request transitions,
 - We are enforcing *agreement*,
 - Dangling states (i.e., unreachable)
- Minimal intervention of the orchestrator (*maximally permissive*)

Example 1 (requests are optional)

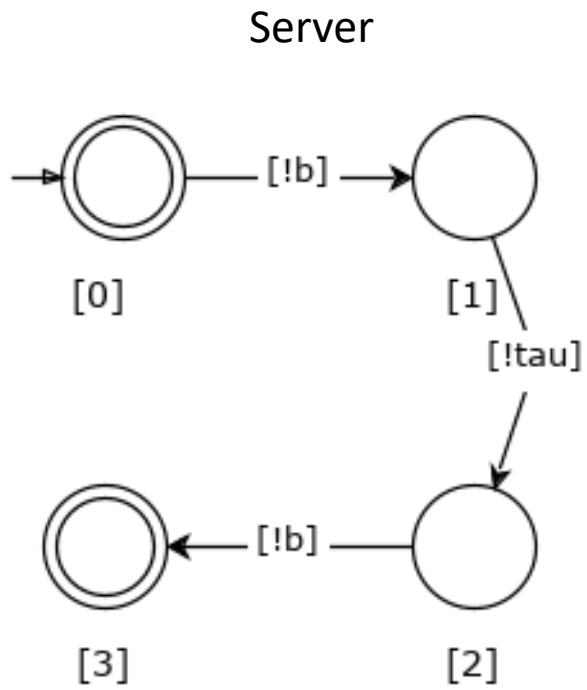
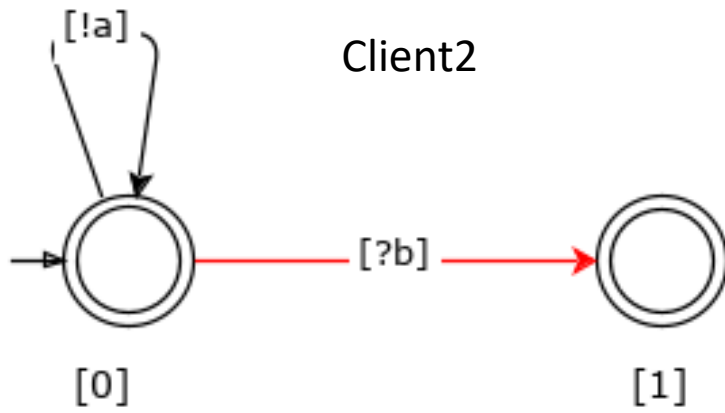


Example 2 (requests are necessary)

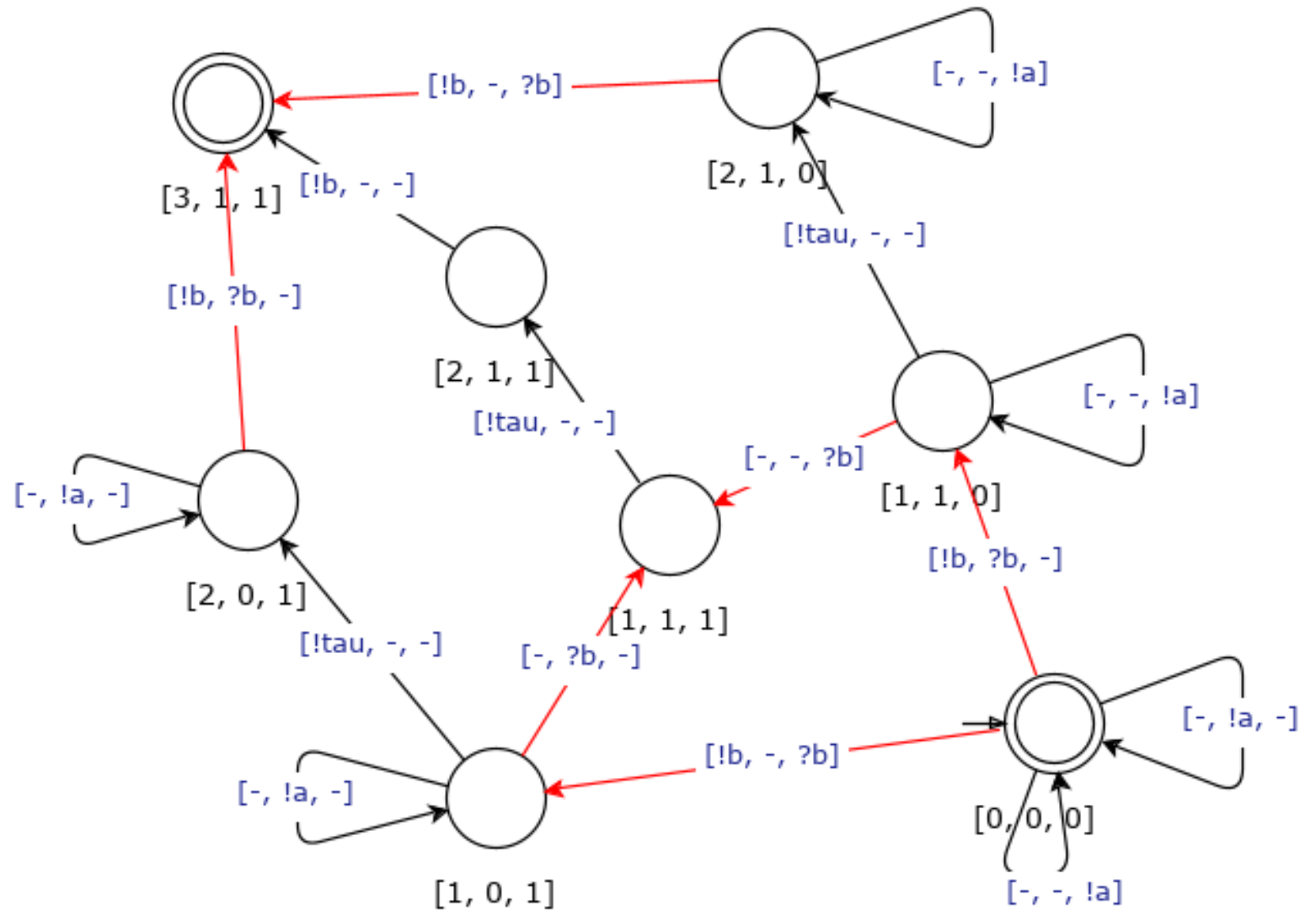
- The necessary request $[?b]$ is *uncontrollable*
- $O(\text{Client2} \times \text{Client2})$ is empty, the initial state is *forbidden!*



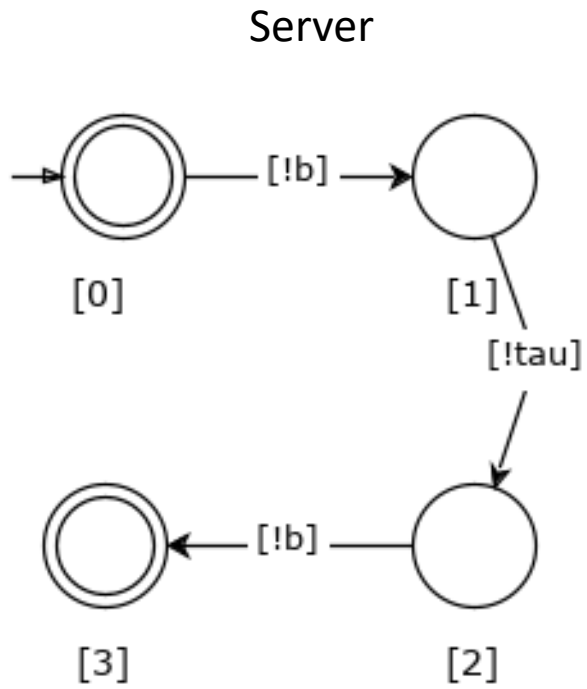
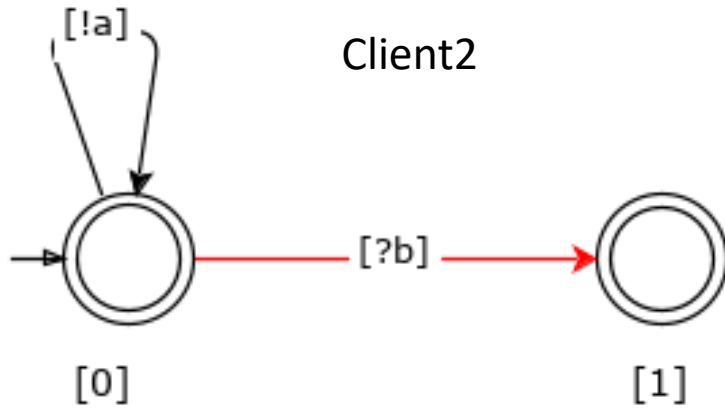
Example 3



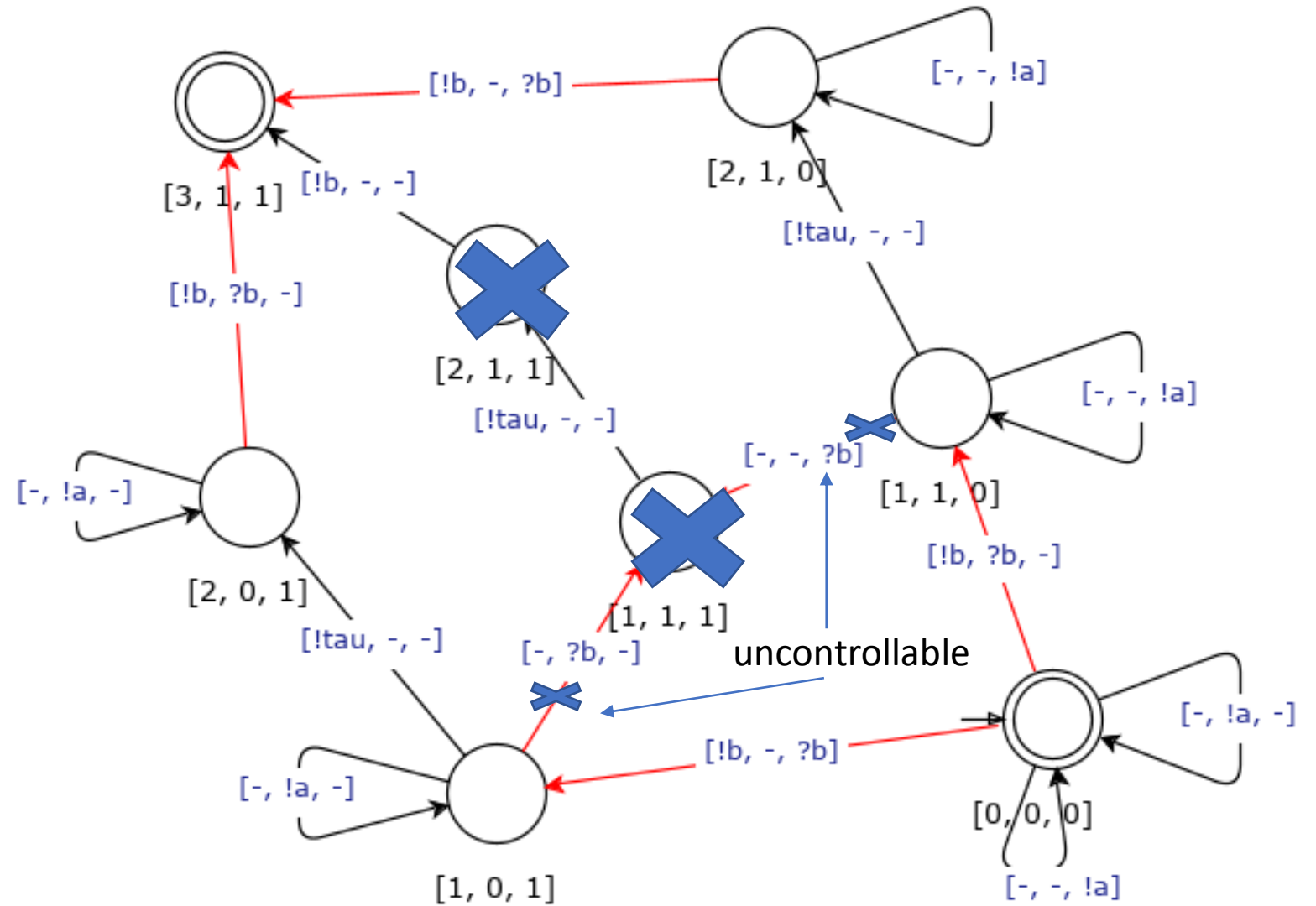
Server x Client2 x Client2



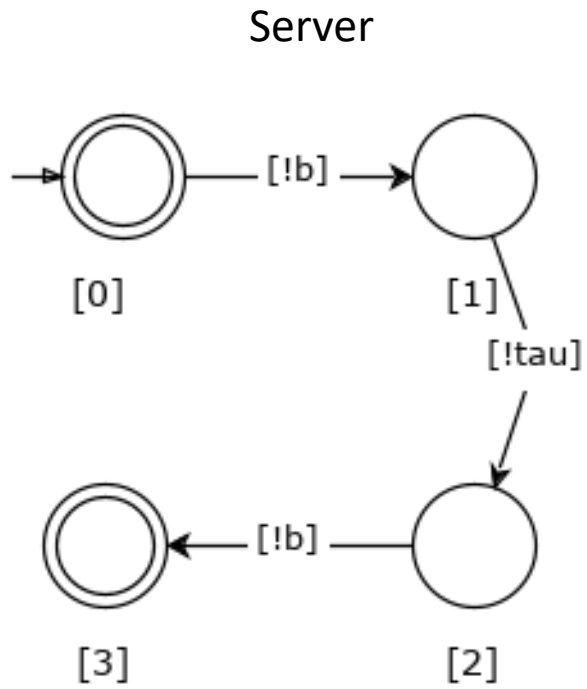
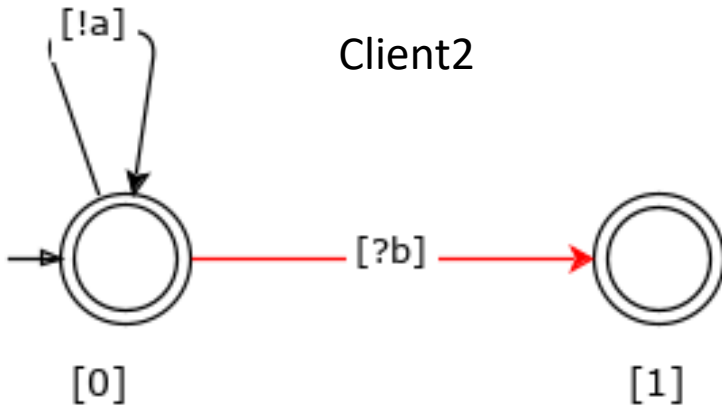
Example 3



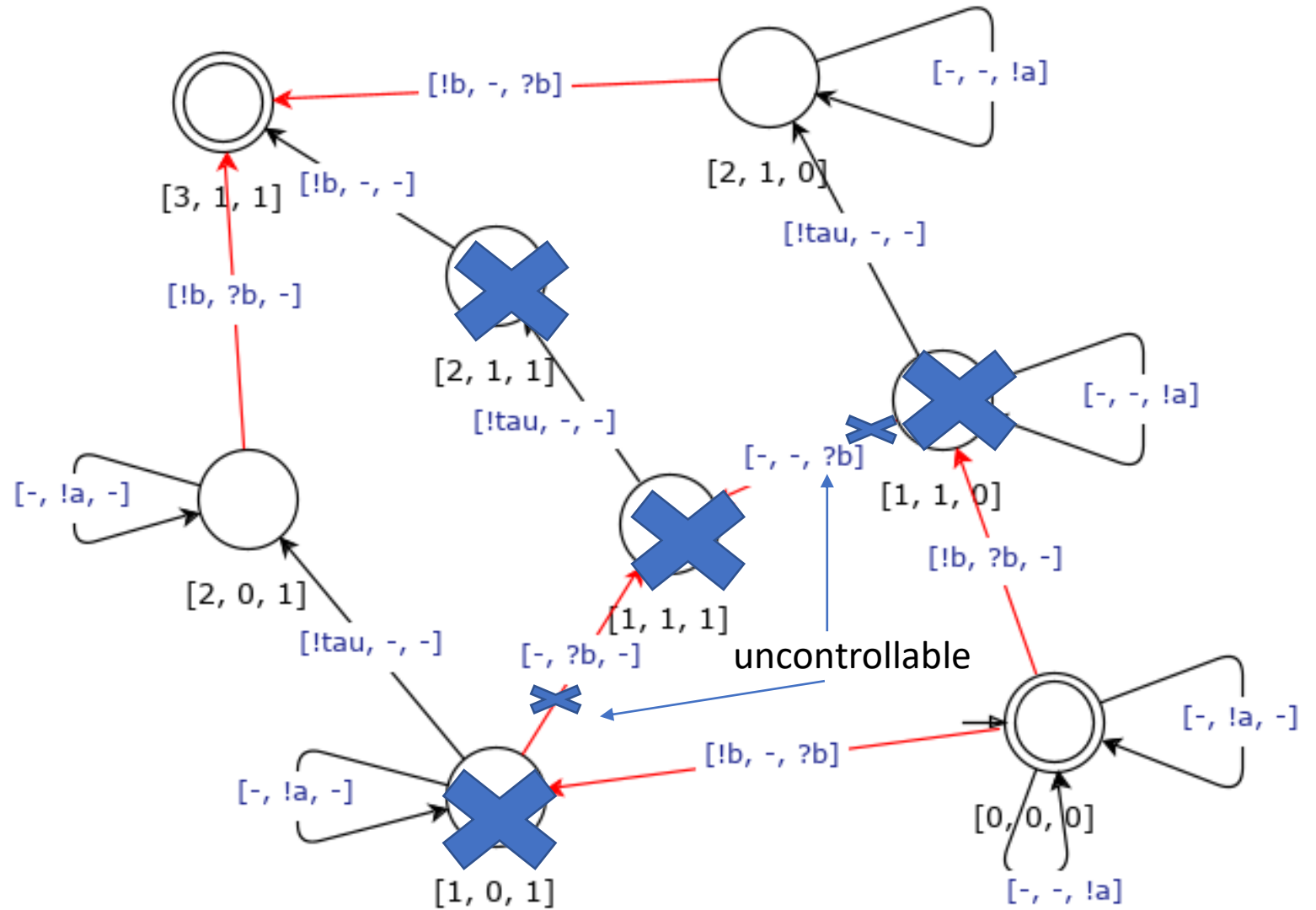
$O(\text{Server} \times \text{Client2} \times \text{Client2})$ first iteration



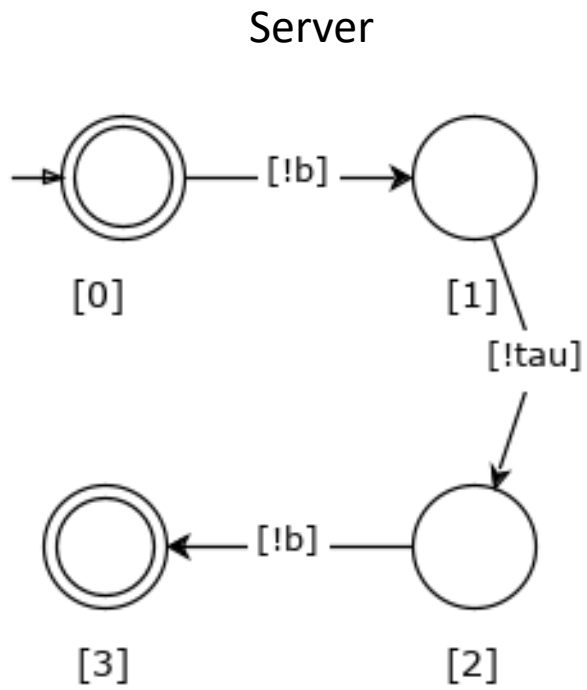
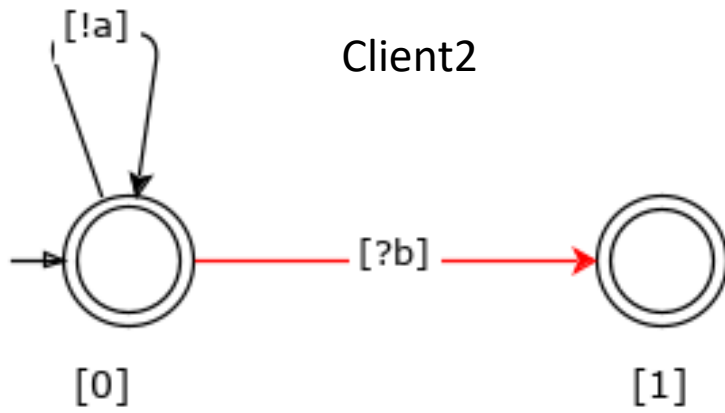
Example 3



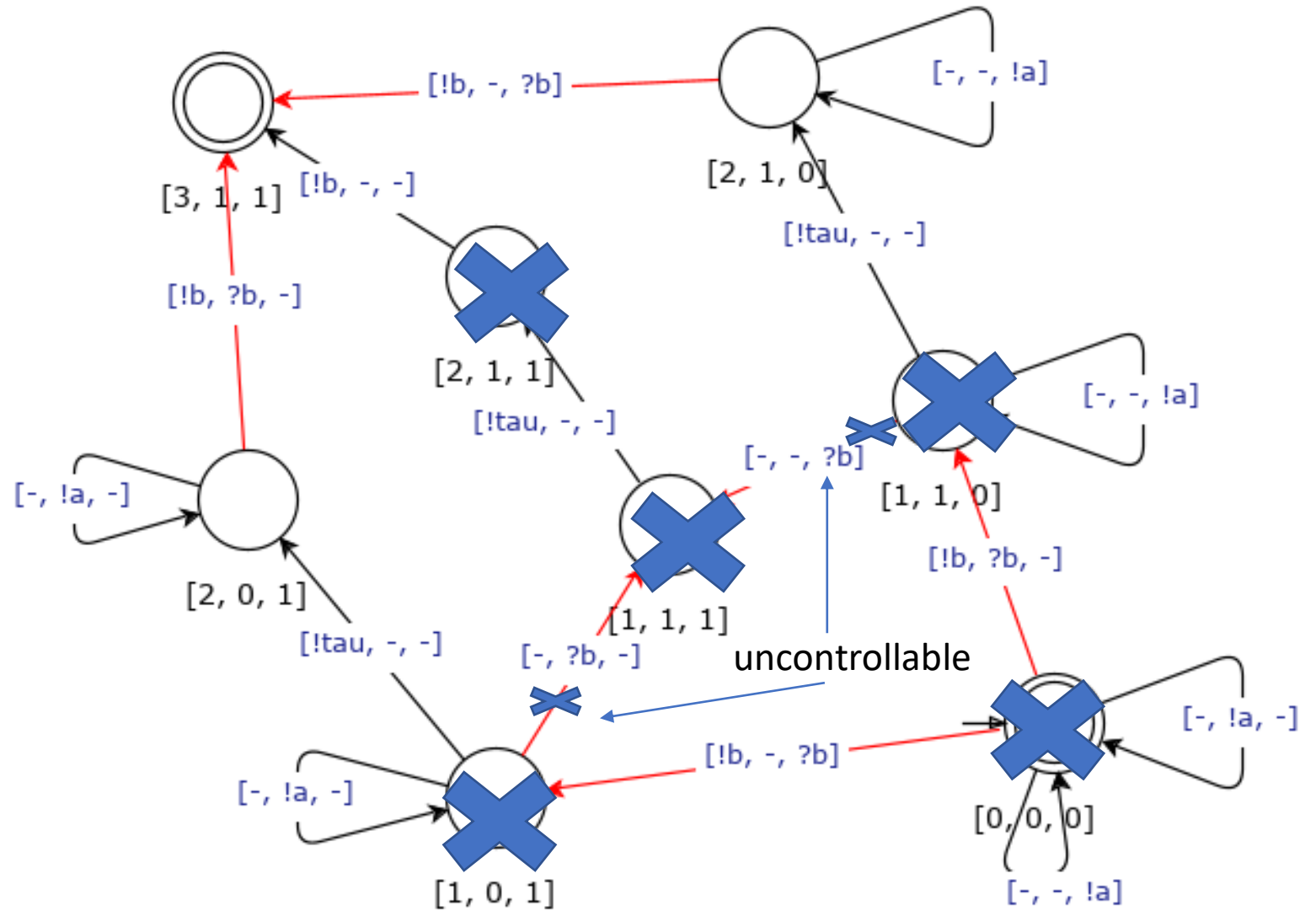
$O(\text{Server} \times \text{Client2} \times \text{Client2})$ first iteration



Example 3

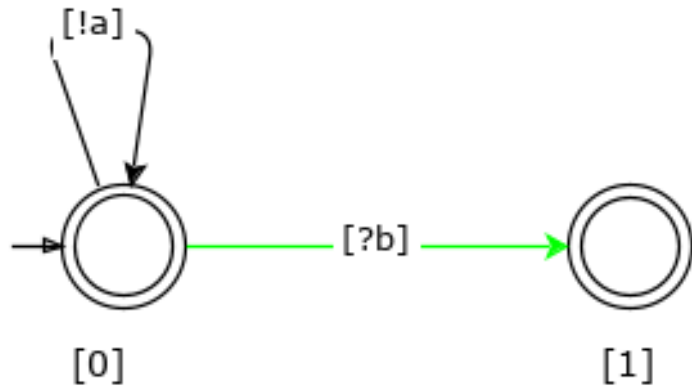


$O(\text{Server} \times \text{Client2} \times \text{Client2})$ fixpoint

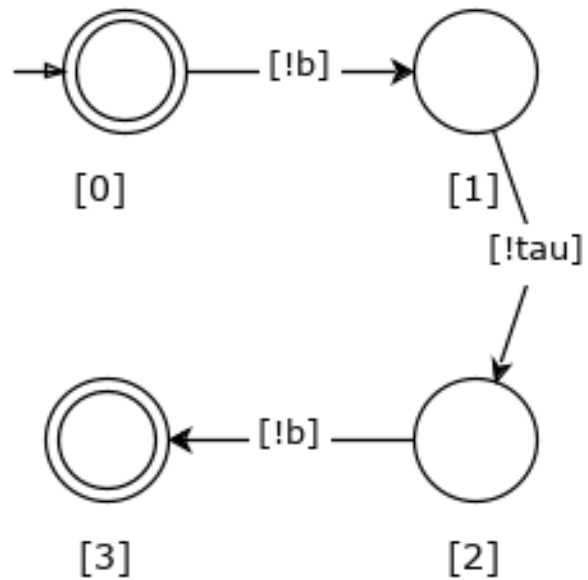


Example 4

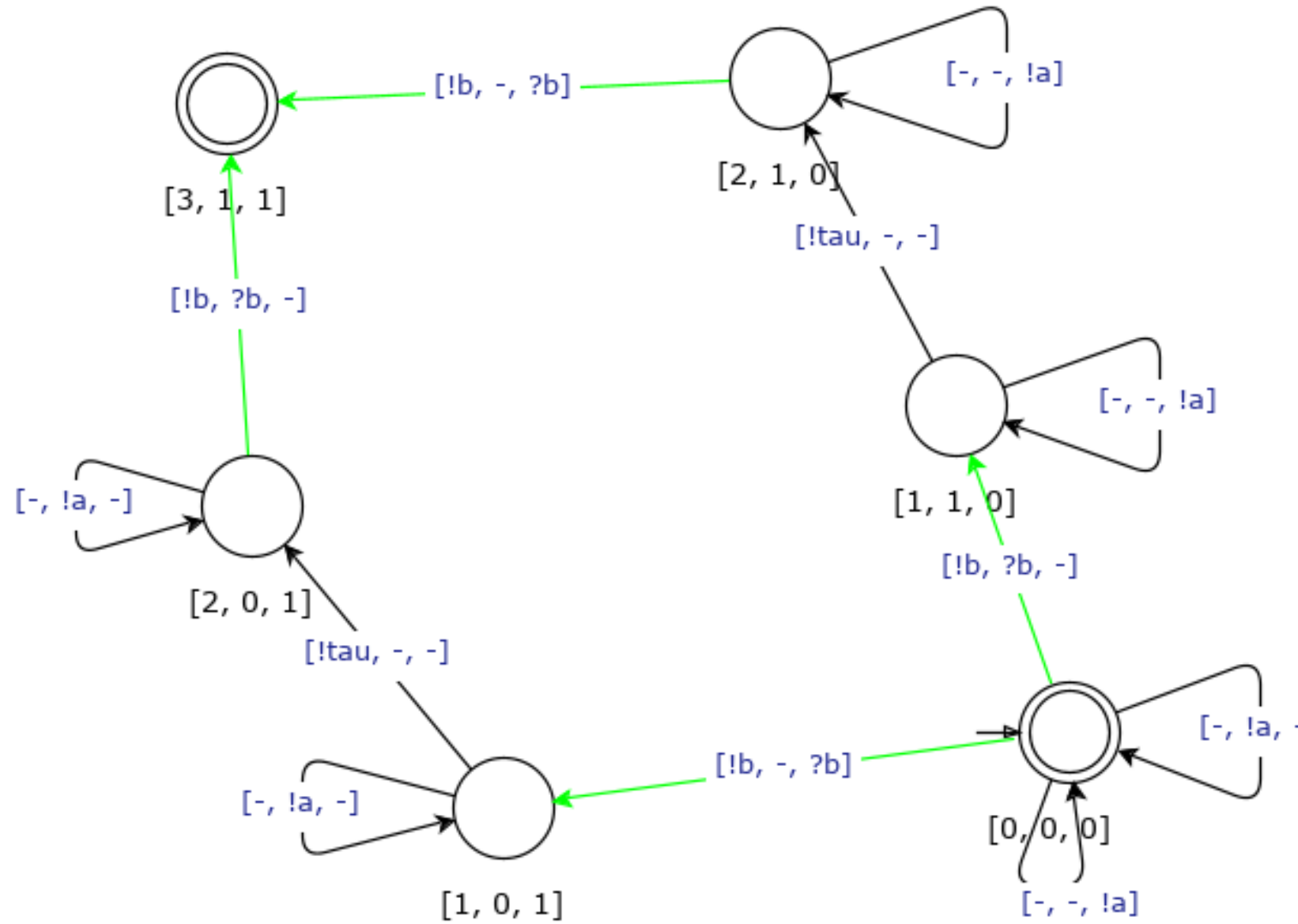
Client2



Server

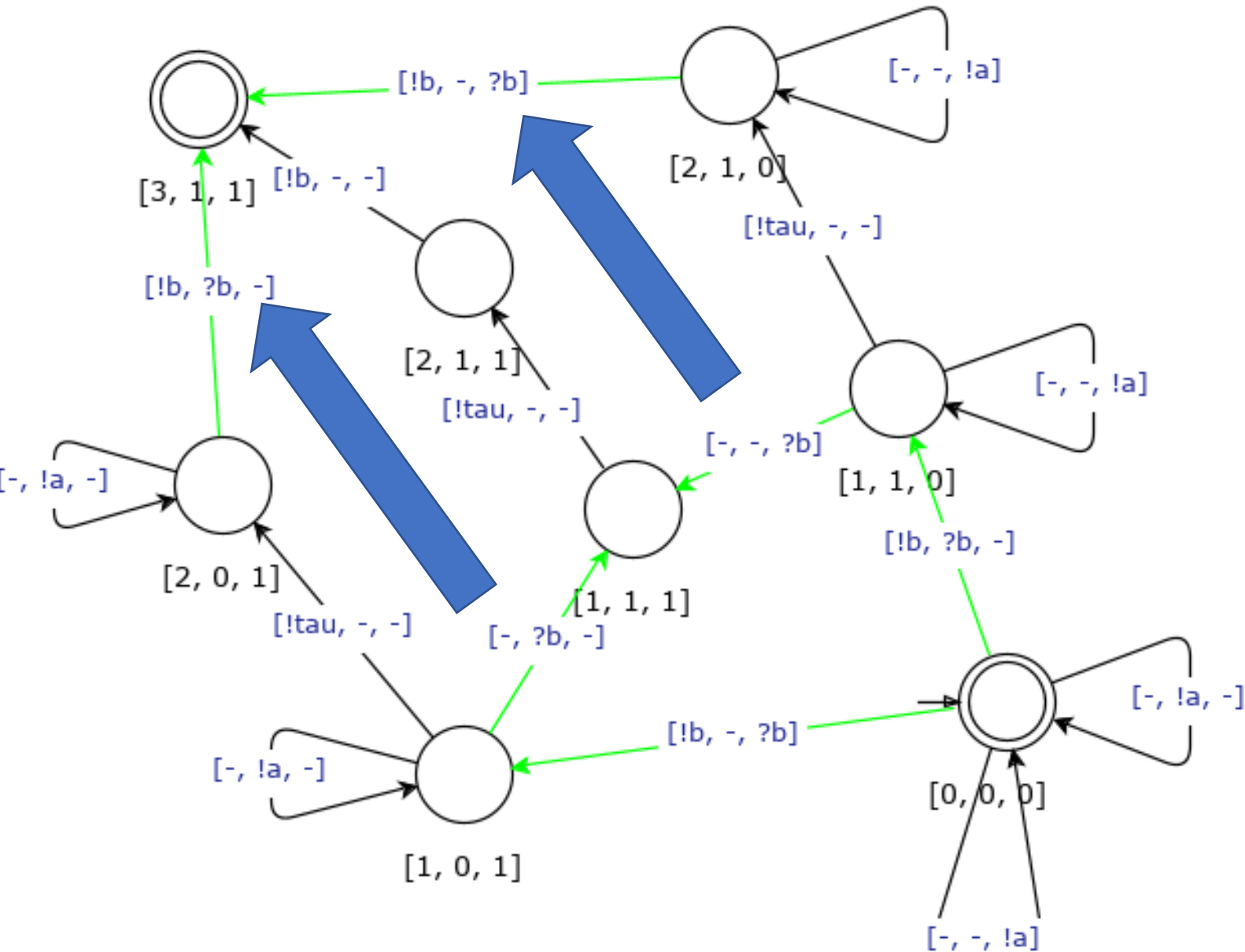


O(Server x Client2 x Client2)



Semi-controllability

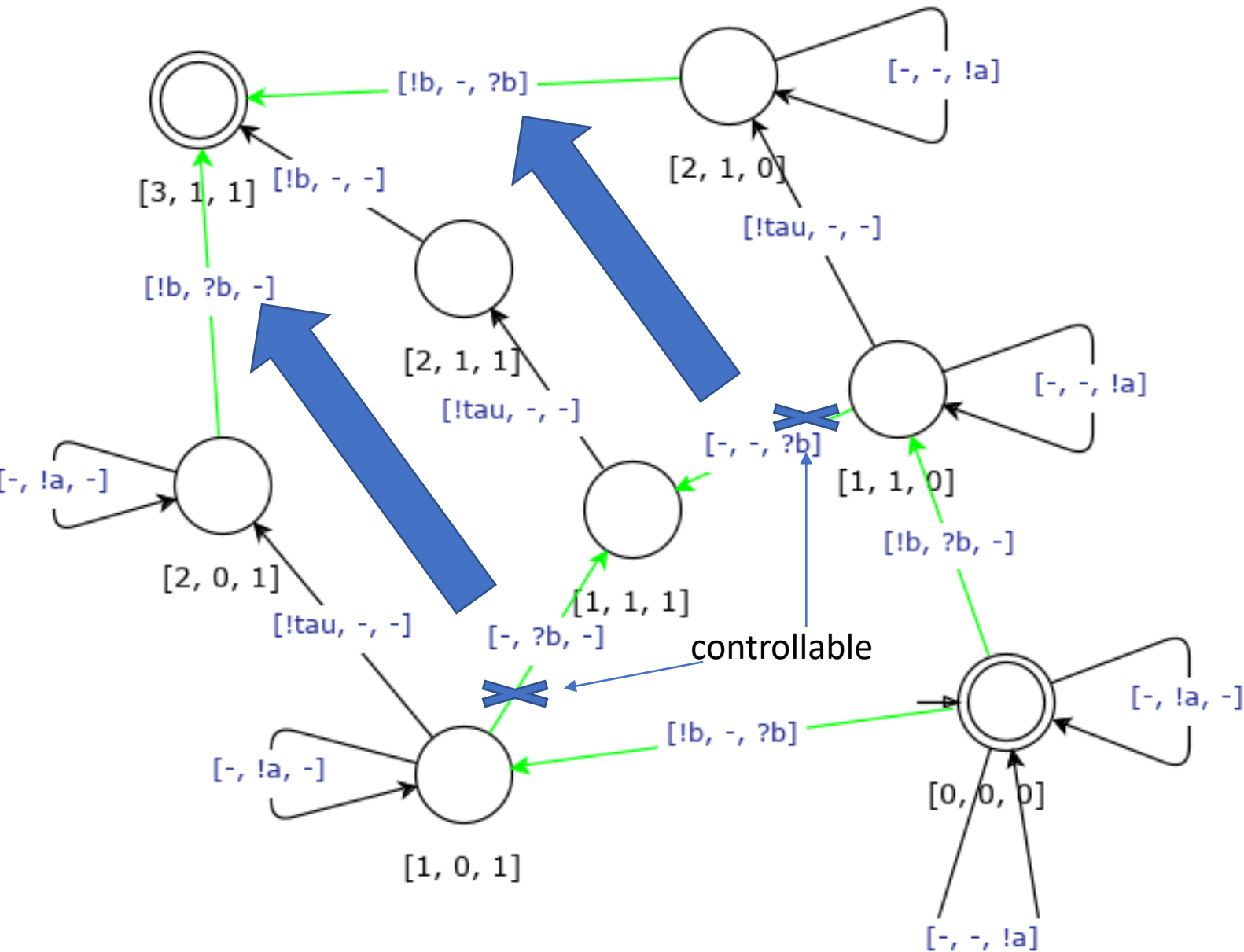
$O(\text{Server} \times \text{Client2} \times \text{Client2})$



Controllable **if** there *exists* another transition matching the same request of the same offerer performed from the same internal state, uncontrollable otherwise.

Semi-controllability

$O(\text{Server} \times \text{Client2} \times \text{Client2})$



Controllable **if** there *exists* another transition matching the same request of the same offerer performed from the same internal state, uncontrollable otherwise.

Contract Automata Runtime Environment

- **CARE** is a middleware recently introduced for realizing applications specified via contract automata
 - It provides an implementation of an *orchestrator*
 - Two aspects to concretise in CARE are
 - Implementation of actions
 - Implementation of choices
 - Abstracted in contract automata
-
- Basile, D. and ter Beek, M.H., 2023, March. A runtime environment for contract automata. In *FM2023*
 - <https://github.com/contractautomataproject> (tools: CATLib, CARE, CATApp)

Example 5

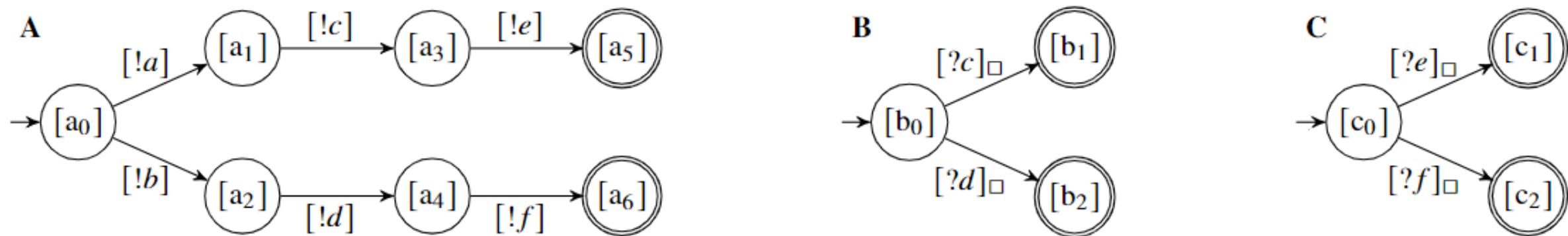


Figure 4: Contracts of **Alice**, **Bob** and **Carl**

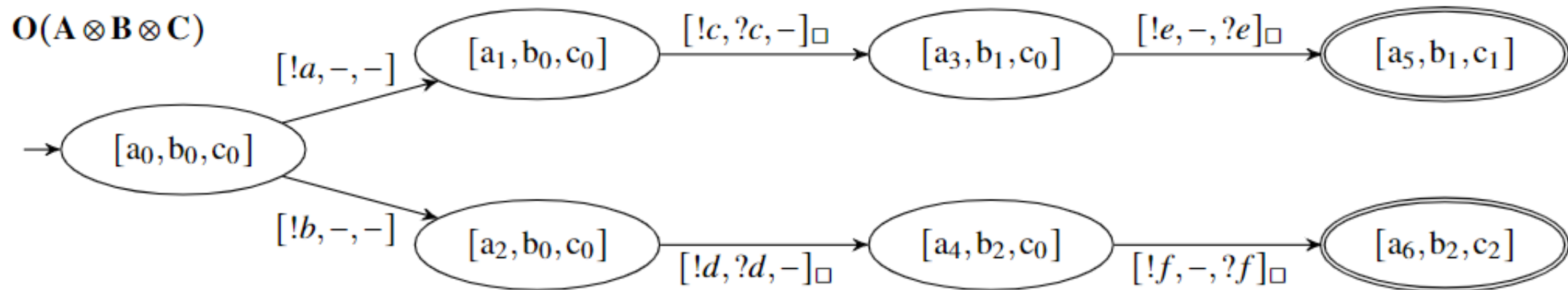
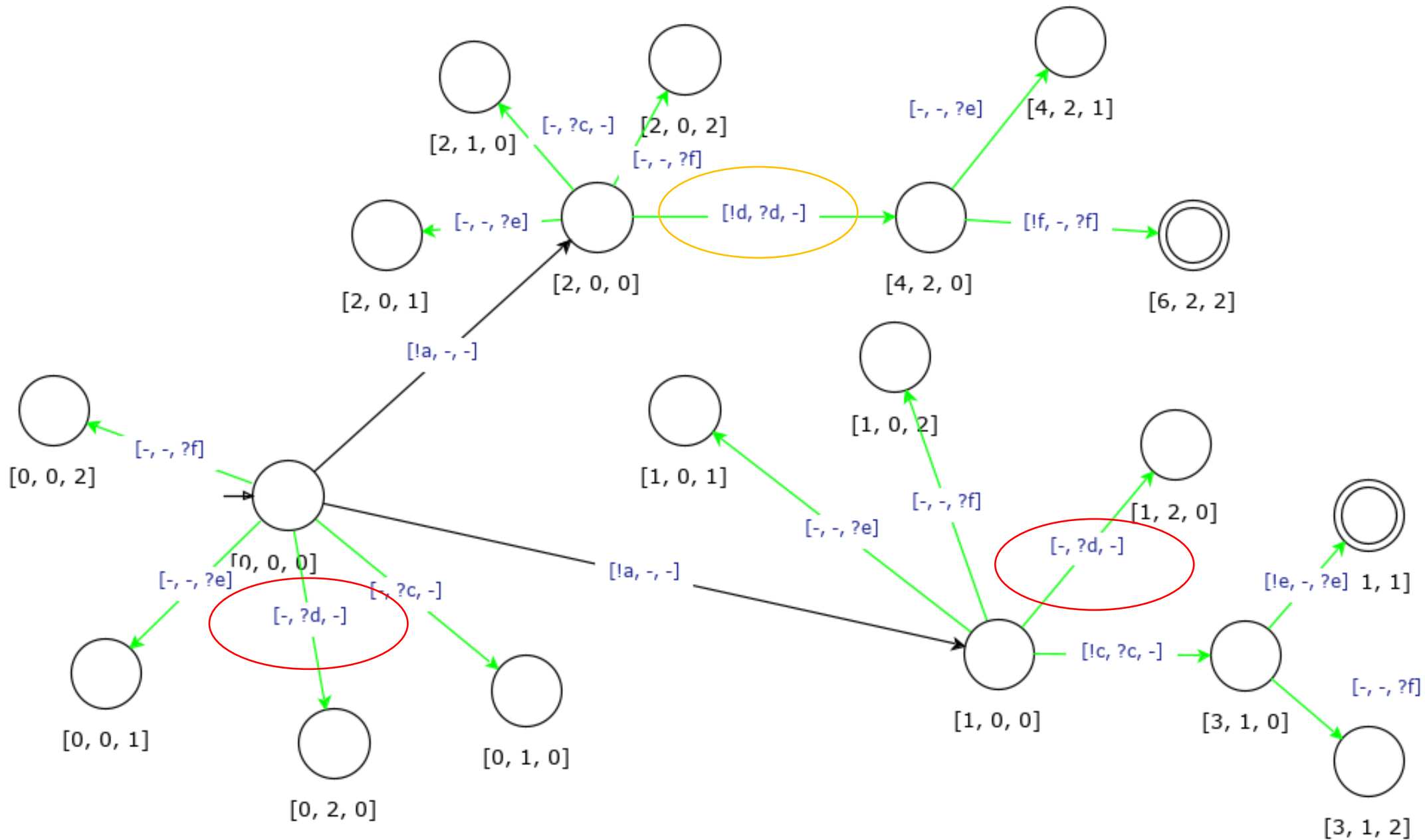


Figure 5: Orchestration $O(A \otimes B \otimes C)$ of **Alice** \otimes **Bob** \otimes **Carl**

Refined Semi-controllability also require that $[2, 0, 0]$ is reachable from $[1, 0, 0]$



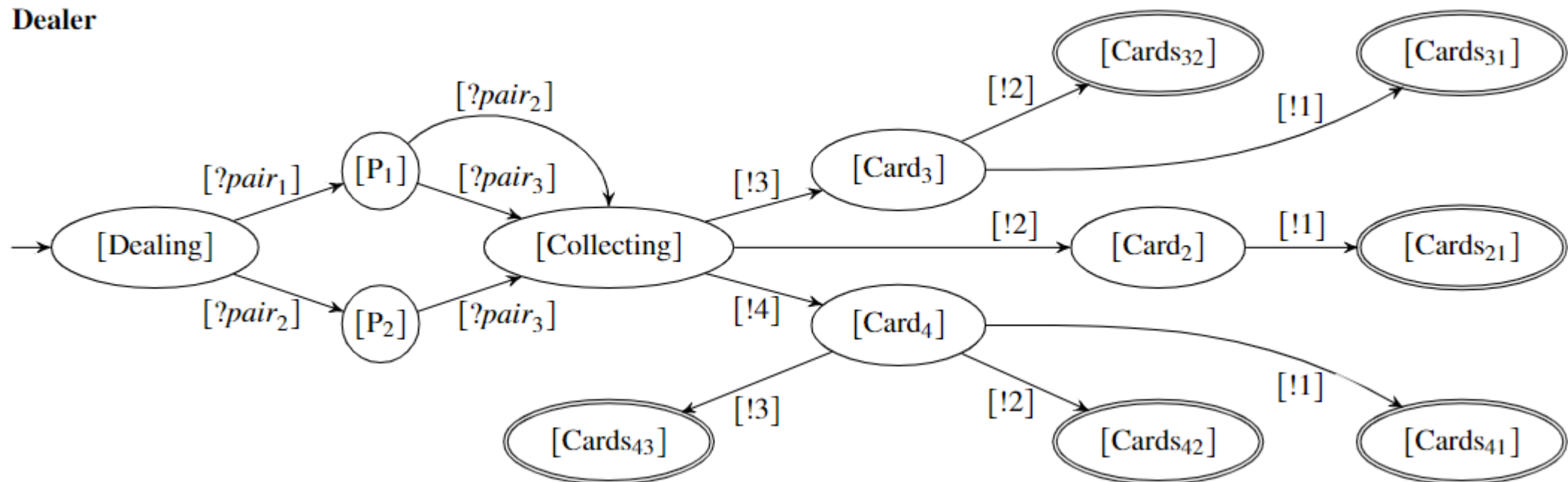
Intuition of necessary service requests

- A service *internally* decides whether to perform a necessary request
- The orchestrator *controls* the scheduling of the requests

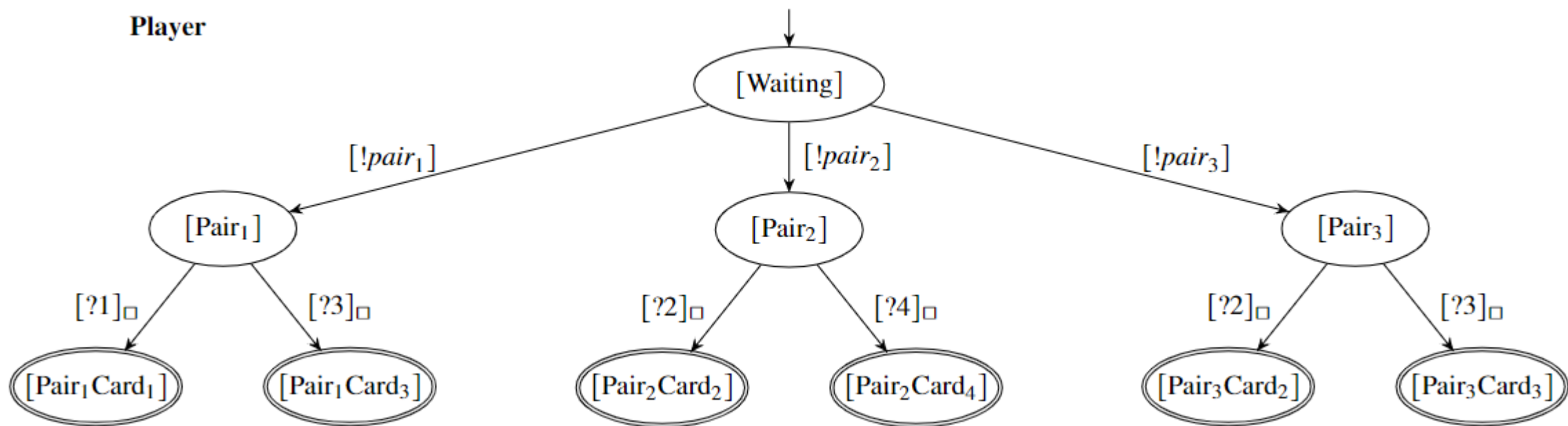
Card dealer example

- Each player receives a pair of cards selected by the dealer from
 - (1,3), (2,4), (2,3)
- Each player picks one of the two cards
- The dealer shall collect the selected cards in descending order
- Strategy of the dealer:
 - Deal in no particular order the pairs (1,3) and (2,4)

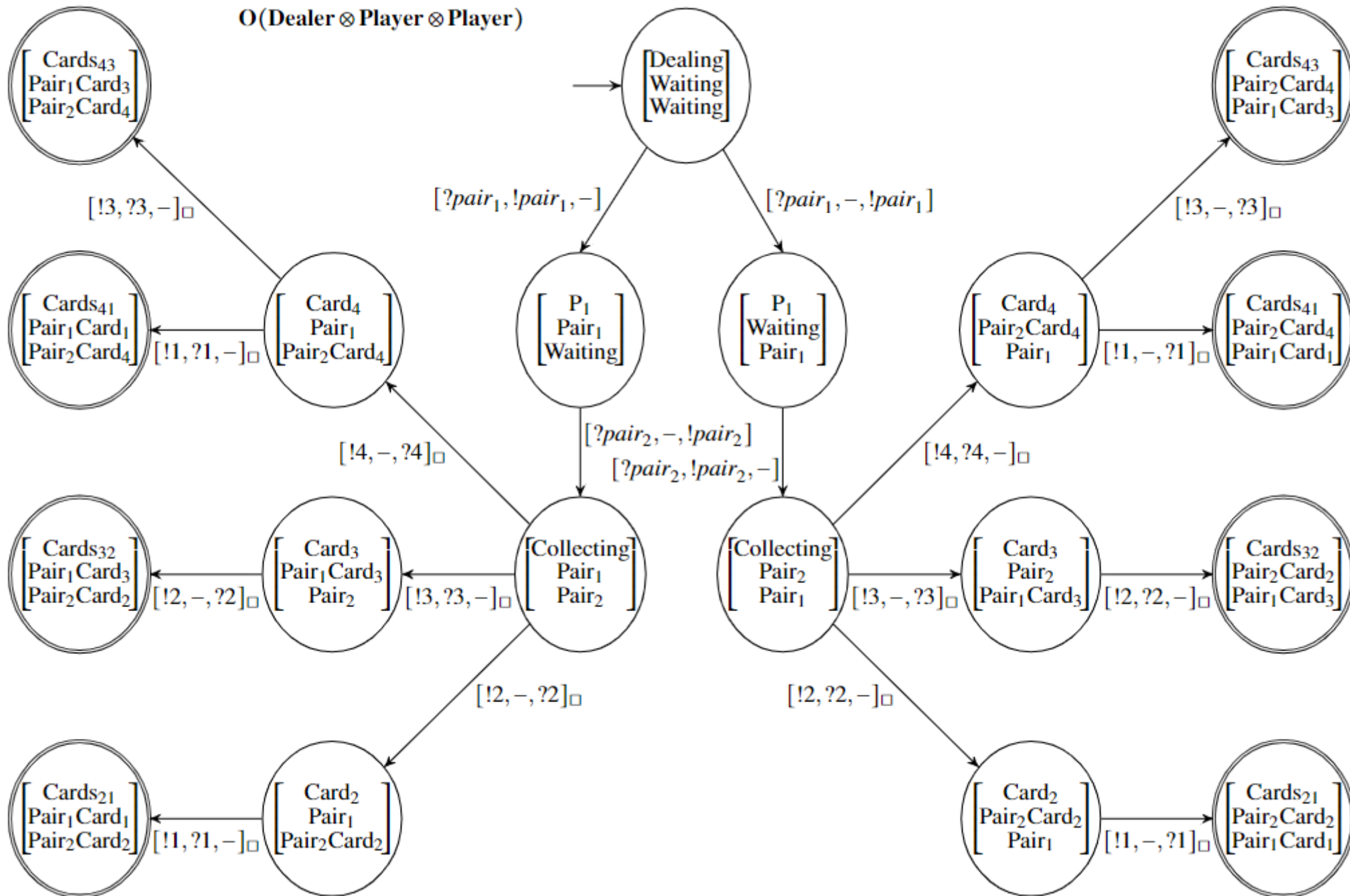
Dealer



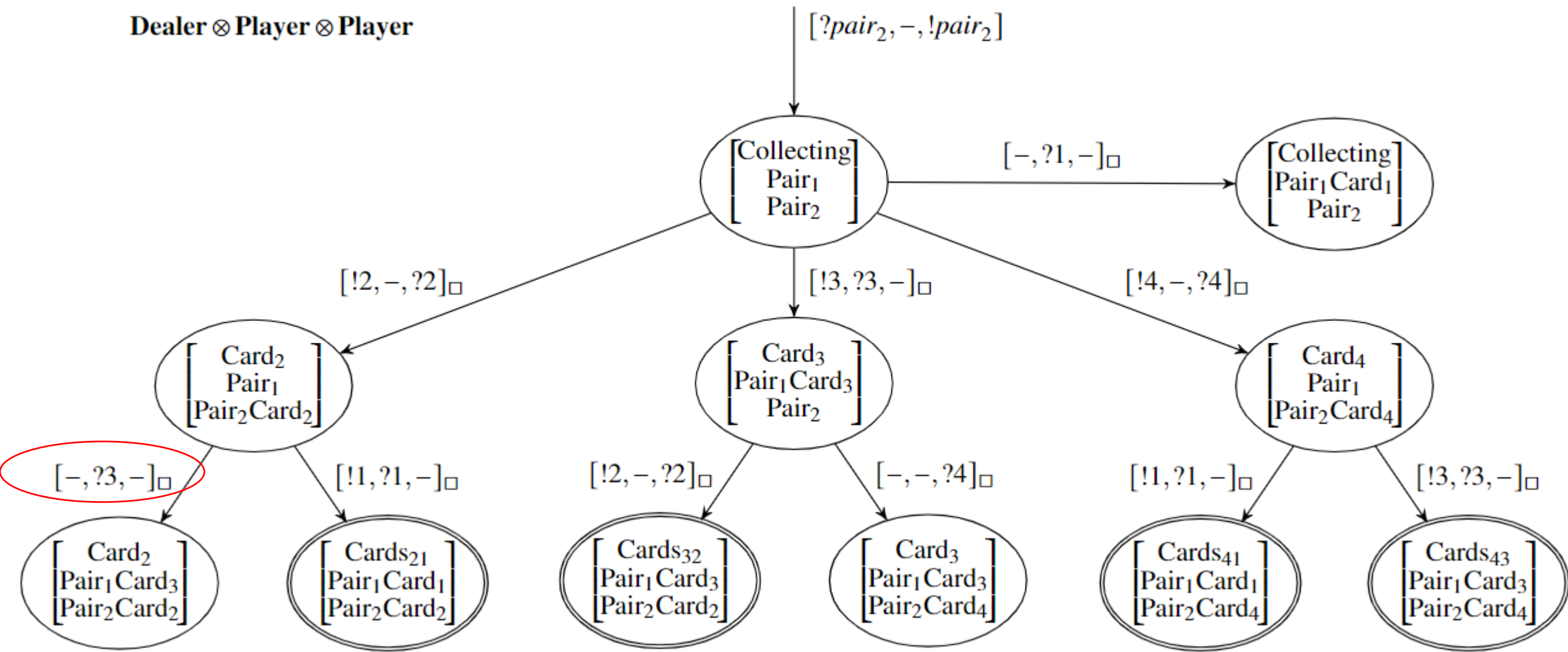
Player



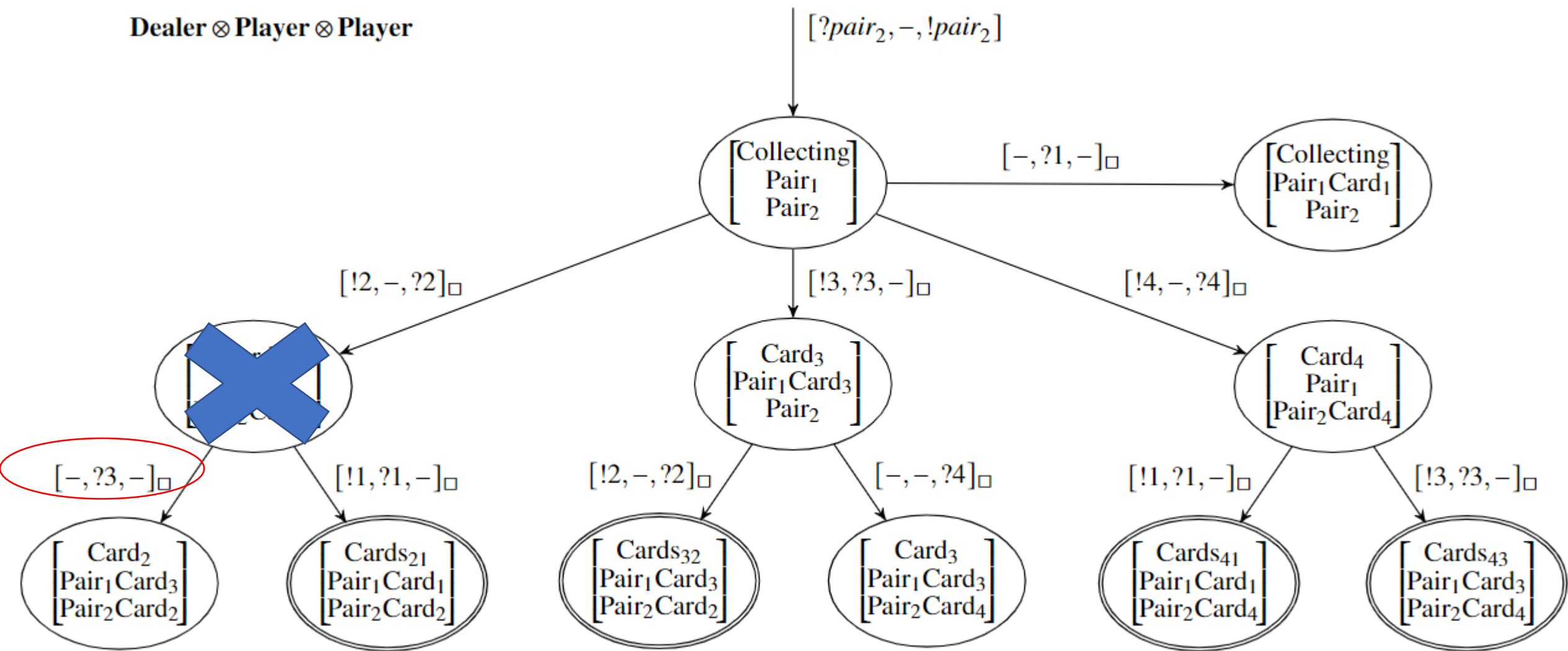
$O(\text{Dealer} \otimes \text{Player} \otimes \text{Player})$



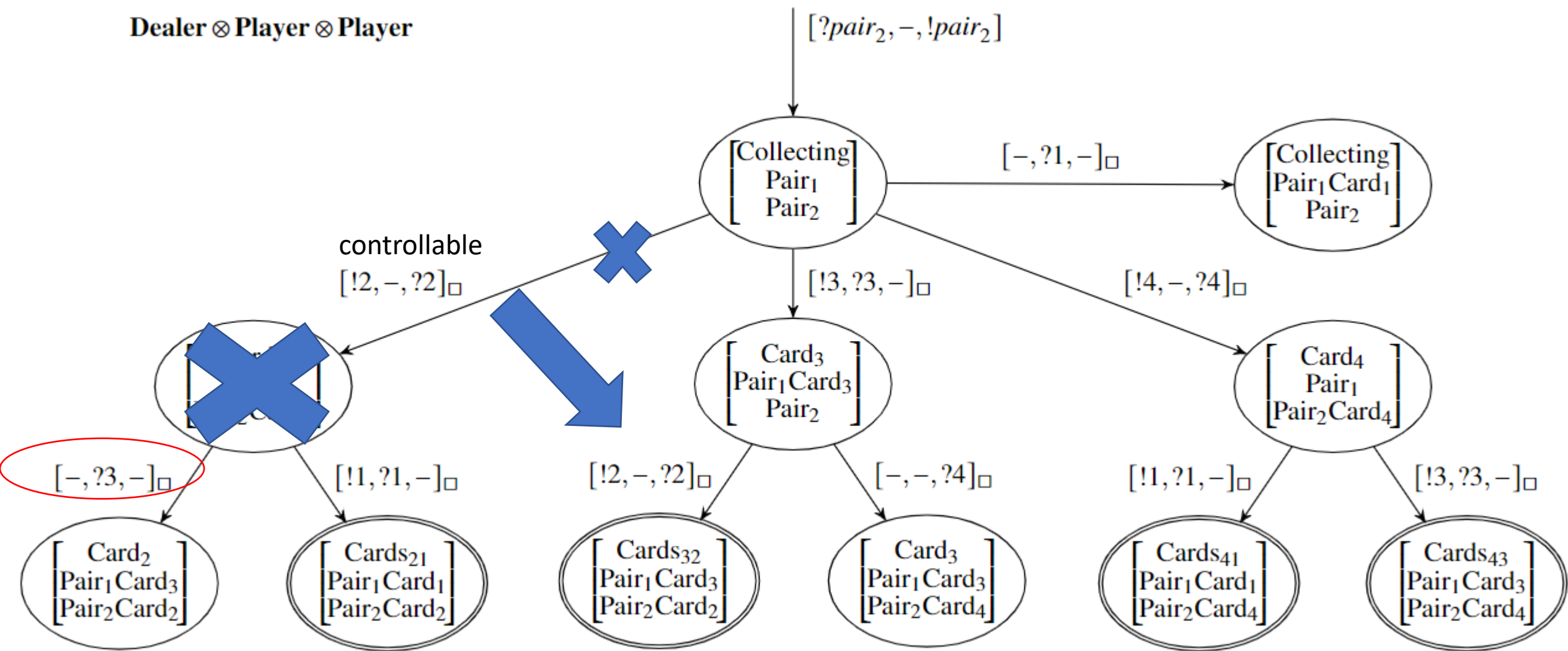
Dealer \otimes Player \otimes Player



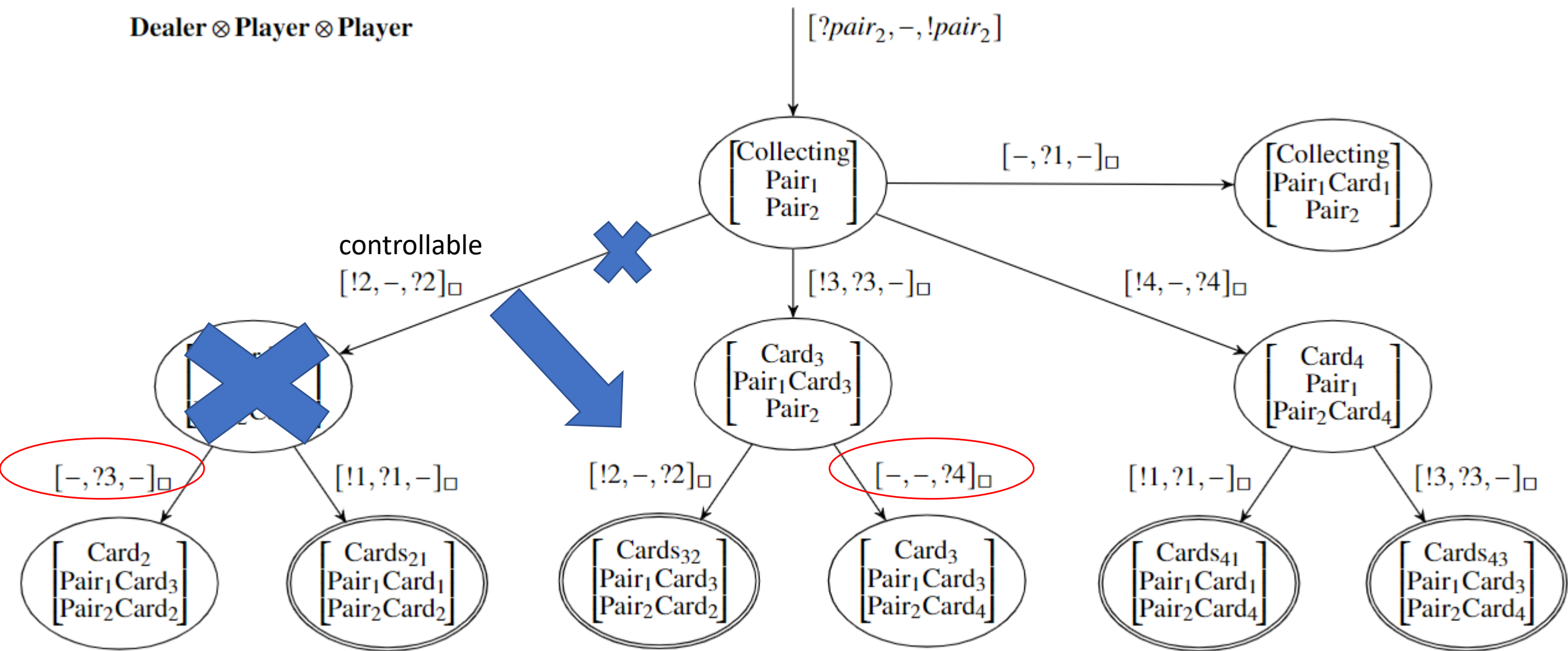
Dealer \otimes Player \otimes Player



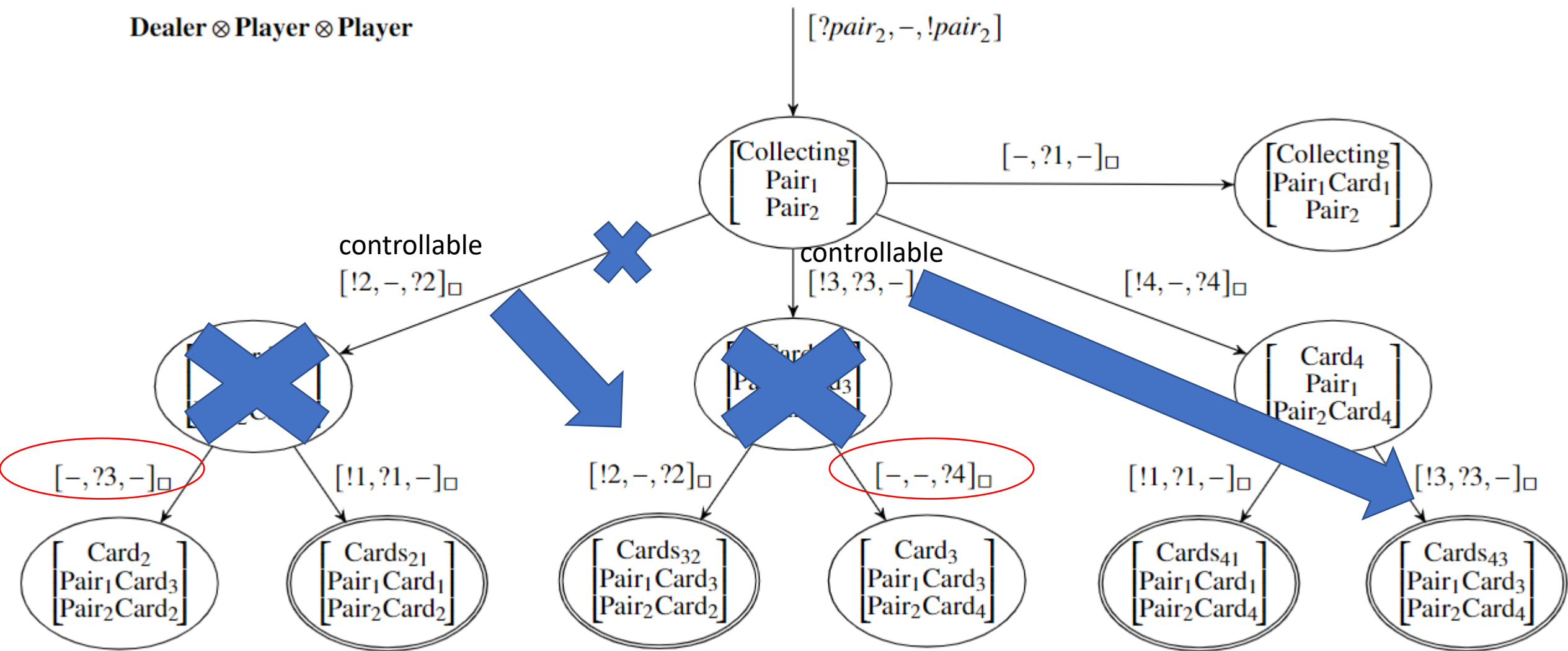
Dealer \otimes Player \otimes Player



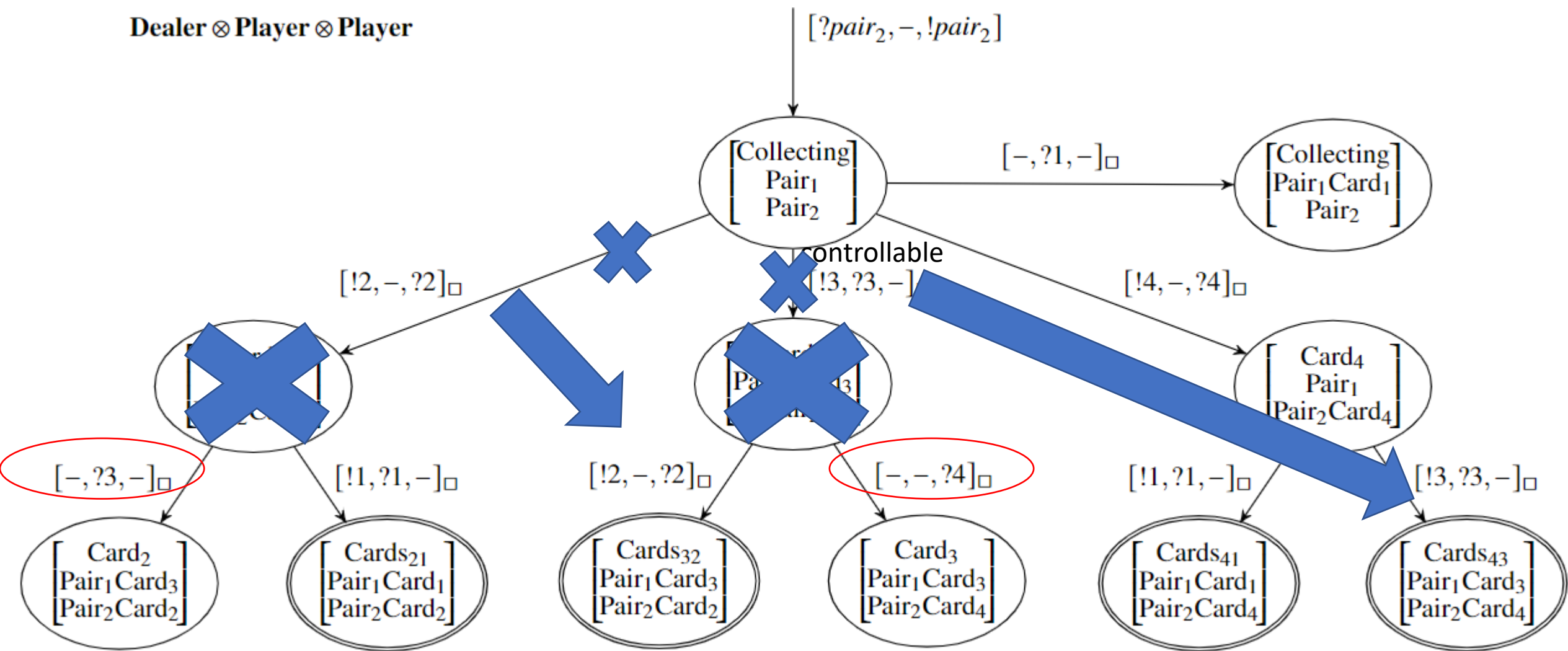
Dealer \otimes Player \otimes Player



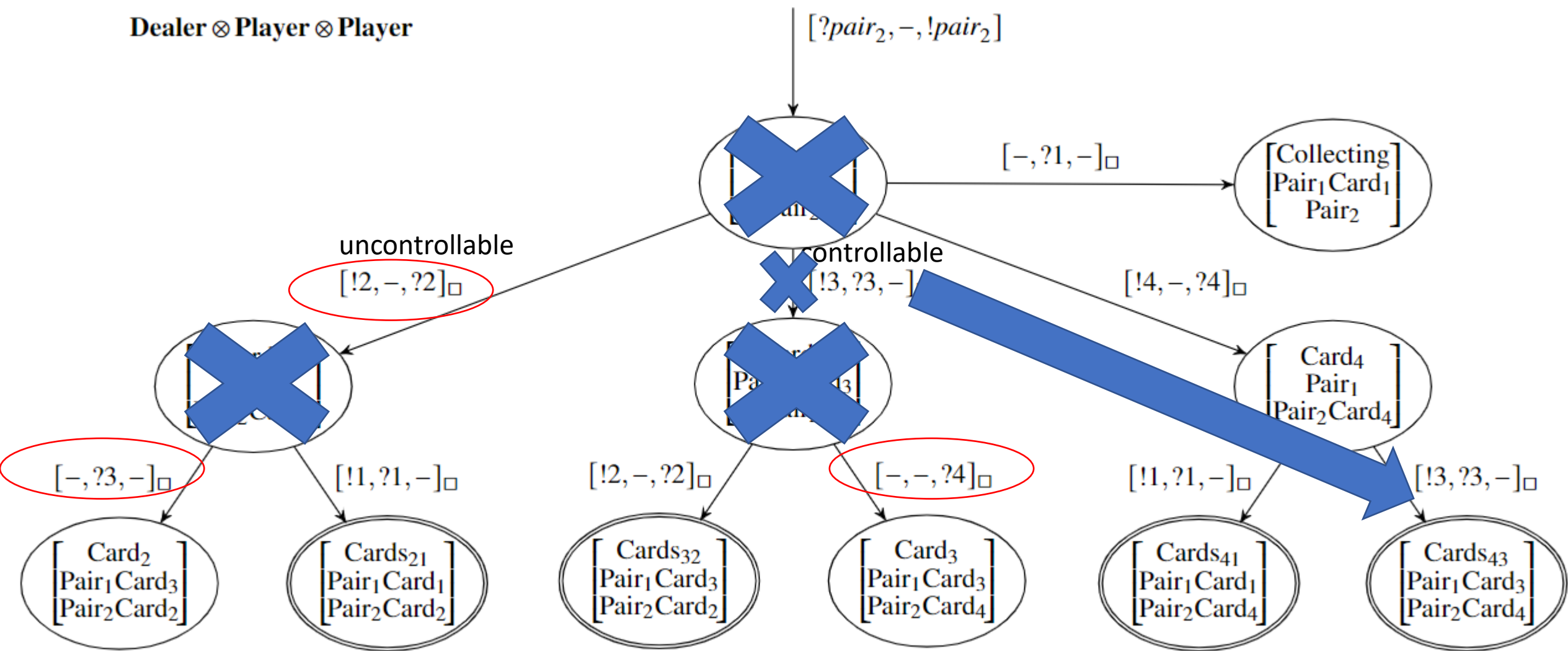
Dealer \otimes Player \otimes Player



Dealer \otimes Player \otimes Player



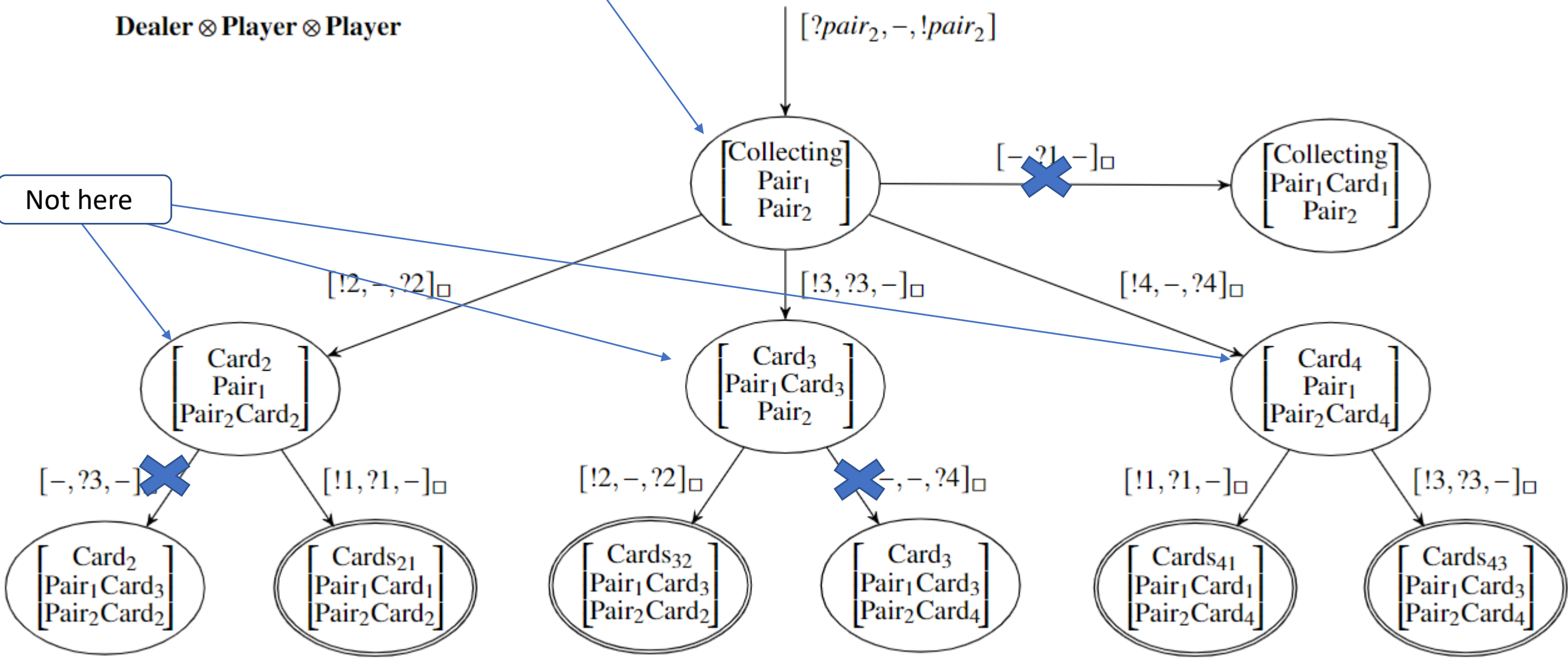
Dealer \otimes Player \otimes Player



The players decide which card to pick in this state

Dealer \otimes Player \otimes Player

Not here



Conclusion

- We have discussed a refinement of the notion of semi-controllability and open challenges in the synthesis of orchestrations in contract automata
- The paper also indicates further challenges and a research roadmap to tackle these challenges effectively

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- We have discussed a refinement of the notion of semi-controllability and open challenges in the synthesis of orchestrations in contract automata
- The paper also indicates further challenges and a research roadmap to tackle these challenges effectively
- **Thanks for your attention!**