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# ROBUST LOGIC PROGRAMMING

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# ROBUST LOGIC PROGRAMMING

## INTEREST OF THIS WORK:

- **IMMEDIATE:** PROGRAMMING TECHNIQUES TO IMPROVE PRODUCTS BASED ON PROLOG
- **GENERAL:** PROLOG AS AN EXAMPLE OF {NON-PROCEDURAL, NON-VON NEUMANN, ARTIFICIALLY INTELLIGENT} {LANGUAGE, PROGRAMMING, COMPUTING}

HOW DOES IT FAIL? HOW CAN FAULTS BE TOLERATED?

## A.I. AND DEPENDABILITY

"A.I." MAY MEAN (PARNAS):

- **APPLICATIONS** THAT ARE NOT YET COMMON AND WELL UNDERSTOOD; HERE DEPENDABILITY PROBLEMS ARE RELATED TO FUZZINESS IN REQUIREMENT OR DIFFICULTIES IN MEASURING PERFORMANCE
  
- **METHODS** LIKE RULE-BASED PROGRAMMING.  
WE ARE INTERESTED IN DEPENDABILITY PROBLEMS SPECIFIC TO THESE METHODS: WE USE PROLOG AS AN EXAMPLE, AND THEN SEE WHICH CONCLUSIONS CAN BE GENERALIZED.

IS PROLOG TYPICAL OF A.I. METHODS? AT LEAST, IT IS PRAISED AND SOLD AS A PROGRAMMING LANGUAGE FOR A.I..

## **PROLOG**

A NON-IMPERATIVE LANGUAGE.

A PROLOG PROGRAM (OR *DATABASE*) IS JUST A DESCRIPTION OF THE PROBLEM.

AN INTERPRETER (A PROBLEM SOLVING PROGRAM) CAN THEN USE THE PROGRAM TO ANSWER QUERIES.

INTERPRETER MAY MEAN: A CONVENTIONAL INTERPRETER, DEDICATED HARDWARE, RULES FOR COMPILING A PROGRAM INTO PROCEDURAL INSTRUCTIONS FOR SEARCHING THE DATABASE

child(jane,jill).  
female(jill).  
female(jane).  
male(james).  
male(arthur).

daughter(X,Y) :- child(X,Y), female(X).  
son(X,Y) :- child(X,Y), male(X).  
mother(X,Y) :- child(Y,X), female(Y).  
father(X,Y) :- child(Y,X), male(Y).  
parent(X,Y) :- mother(X,Y).  
parent(X,Y) :- father(X,Y).

EXAMPLES OF QUERIES:

?- male(jill).

no

?- daughter(X,Y).

X = jane, Y = jill ;

no

?-

## PROS AND CONS OF PROLOG.

### THEORETICAL ADVANTAGES:

- **SIMPLICITY** - THE PROGRAMMER CAN DISREGARD HOW SOLUTIONS ARE OBTAINED AND CONCENTRATE ON WHAT THEY SHOULD BE;
- **EFFICIENCY** - THE INTERPRETER CAN BE MADE FASTER (E.G. BY PARALLELISM) WITHOUT CHANGING THE PROGRAMS;
- **FLEXIBILITY** - A DATA BASE CAN BE QUERIED TO OBTAIN ANY RESPONSE THAT IS DEDUCTIBLE FROM IT.

## **DISADVANTAGES:**

- **INEFFICIENCY** A GENERAL-PURPOSE INTERPRETER IS GENERALLY INEFFICIENT; A PROGRAM THAT ANSWERS SOME QUERIES EFFICIENTLY MAY BE VERY INEFFICIENT WITH OTHER QUERIES.

THIS REQUIRES: A) THE PROGRAMMER TO PAY MUCH ATTENTION TO HOW SOLUTIONS ARE OBTAINED; B) THE IMPLEMENTORS TO ADD A LOT OF PROCEDURAL AND METALOGICAL CONSTRUCTS.

- **PROGRAMS USED IN UNFORESEEN WAYS MAY FAIL IN UNFORESEEN WAYS**

## **EXAMPLES OF NON-LOGICAL BUILT-IN PREDICATES:**

**CUT:** FREEZES THE CHOICES MADE FOR SOME VARIABLES SO THAT BACKTRACKING CANNOT CHANGE THEM

**ASSERT:** ADDS A CLAUSE TO THE DATABASE

**RETRACT:** REMOVES A CLAUSE FROM THE DATABASE

**VAR, NONVAR:** TRUE IF A TERM IS (IS NOT) AN UNINSTANTIATED VARIABLE WHEN THE PREDICATE IS ENCOUNTERED DURING EXECUTION

ETC.



## **RELIABILITY PROBLEMS WITH PROLOG:**

### **ACCIDENTAL: ARISING FROM YOUTH:**

- MANY PROLOG ENVIRONMENTS ARE TOYS
- IDEAL PURITY IS CONSIDERED MORE THAN PRACTICAL NECESSITIES

**LACK OF:** TYPES, MODULES, NAME SCOPE LIMITATIONS, EXCEPTION HANDLING. MOST TYPING ERRORS DON'T CAUSE ANY SYNTACTIC ERROR: THEY JUST CREATE A LEGAL WRONG PROGRAM. (RECENT PROLOG IMPLEMENTATIONS INCLUDE MANY OF THESE FEATURES. FURTHER IMPROVEMENTS ARE LIKELY WITH NEWER DEVELOPMENT ENVIRONMENTS AND/OR LANGUAGES)

**LACK OF STANDARDS:** DIFFERENT IMPLEMENTATIONS HAVE DIFFERENT SEMANTICS AND INCONSISTENT LIBRARIES

## PROBLEMS INHERENT IN:

- THE SEQUENTIAL SEARCH STRATEGY OF ORDINARY PROLOG ?
- THE PROGRAMMING STYLE BASED ON FIRST-ORDER LOGIC?
- NON-PROCEDURAL PROGRAMMING?
- .....

?

## **A POSSIBLE OBJECTION:**

LOGICAL PROGRAMMING DOES NOT NEED  
FAULT-TOLERANCE, BECAUSE:

- A. PROLOG IS LOGIC: YOU CAN PROVE  
CORRECTNESS IF YOU CARE TO;
  
- B. A.I. PROGRAMS ARE NATURALLY ROBUST:  
IF THEY CANNOT FIND A SOLUTION AT  
FIRST, THEY WILL KEEP SEARCHING AND  
FIND ONE IF AT ALL POSSIBLE.

## ANSWERS:

### A. PROLOG PROGRAMS ARE NOT PURE LOGIC.

- THE EXECUTION IS CONTROLLED BY SUCH THINGS AS THE ORDERING OF CLAUSES.
- PROGRAMS THAT ARE CORRECT DESCRIPTIONS OF PROBLEMS AND FAIL TO PRODUCE SOLUTIONS ARE COMMONPLACE.
- SUCCESSFUL PROLOGS HAVE LOTS OF PROCEDURAL OR META-LOGIC PREDICATES, WHOSE MEANING DEPENDS ON DETAILS OF THE WORKINGS OF THE INTERPRETER.
- IT HAS BEEN SHOWN THAT MANY PROLOG PROGRAMS ARE MEANT TO EXECUTE LIKE IMPERATIVE PROGRAMS.
- WHAT ABOUT FAULTS IN THE SUPPORT (HW PLUS OS PLUS INTERPRETER)?

### B. WE WANT CORRECT SOLUTIONS, NOT JUST ANY SOLUTION.

REDUNDANCY IN LOGIC PROGRAMS CAN HAPPEN BY CHANCE AND EVEN BE HARMLESS, BUT USEFUL REDUNDANCY NEEDS PLANNING.

## SOFTWARE FAULT-TOLERANCE:

USING EXTRA INFORMATION ABOUT THE PROBLEM BESIDE THE INFORMATION NECESSARY TO FIND A SOLUTION.

A METHOD IS DEFINED BY:

- WHICH INFORMATION IS PROVIDED (NECESSARY CONDITIONS ABOUT THE SOLUTION, ALTERNATIVE WAYS OF COMPUTING THE SOLUTION)
- HOW THE REDUNDANT INFORMATION IS USED TO OBTAIN A MORE RELIABLE SOLUTION, OR TO SIGNAL AN ERROR

REDUNDANT INFORMATION IS EASILY EXPRESSED IN PROLOG.

EXAMPLES OF FAULT TOLERANCE IN PROLOG PROGRAMS

ASSERTIONS (PRE/POST CONDITIONS)

ASSERTIONS WITH BACKWARD RECOVERY

DIVERSITY WITH ADJUDICATION

```
/* this combination first uses <constructive  
predicates (X,Y,...)> to find a solution, then  
checks it by assertion(X,Y,...) */
```

```
statement(X,Y,...) :- <constructive predicates  
                  (X,Y,...)>, assertion(X,Y,...).  
assertion(X,Y,...) :- nonvar(X),  
                      nonvar(Y),...,<appropriate  
                      predicates>, !.
```

```
/* notice the nonvar() and the cut */
```

```
/* by adding this, we can treat a failed  
assertion explicitly:*/
```

```
assertion(X,Y,...) :- <series of predicates to  
                      handle exception>.
```

```
/* assertions on preconditions */  
statement(X,Y,...) :- assertion(X,Y,...),  
                      <constructive predicates (X,Y,...)>.  
assertion(X,Y,...) :- nonvar(...),...,  
                      <other predicates>.
```

```
/* if we can code two diverse ways to satisfy a  
goal, and want to use backtracking as  
backward recovery */  
recoverable_goal(X,Y,..) :- procedure1(X,Y,..),  
    assertion(X,Y,..).  
recoverable_goal(X,Y,..) :-  
    procedure2(X,Y,..), assertion(X,Y,..).
```

```
/* if we can code two diverse ways to satisfy a goal, and want static redundancy with comparison */
```

```
duplex_goal(In1, In2,..,Out1,Out2,..) :-  
    procedure1(In1,In2,..,Out1,Out2,..),  
    procedure2(In1,In2,..Out1',Out2',..),  
    consistency(Out1,Out2,..,Out1',  
    Out2',..).
```

```
consistency(Out1,Out2,..,Out1',Out2',..) :-  
    <predicates(Out1,Out2,..,Out1',Out2'  
    ',..)>.
```

```
/* we can deal with mismatches by */
```

```
consistency(Out1,Out2,..,Out1',Out2',..) :-  
    <predicates(Out1,Out2,..,Out1',Out2'  
    ,..)>, !.
```

```
consistency(Out1,Out2,..,Out1',Out2',..) :-  
    <exception handling predicates>.
```

```
/* serious problems arise with backtracking and/or multiple solutions */
```



```
/* in general for N diverse implementations */  
Nuplex_goal(In1, In2,...,Out1,Out2,..) :-  
    procedure1(In1, In2,...,Out11, Out21,  
    ..), procedure2(In1,  
    In2,...,Out12,Out22,..), .....,  
    procedureN(In1, In2,...,Out1N,  
    Out2N, ..),  
    adjudicator(Out11,Out21,...,Out1N,  
    Out2N,..).
```

**/\* This required all N procedures to succeed.  
Instead, we may want threshold voting, which  
is cumbersome without explicit non-  
determinism \*/**

## **ASPECTS THAT NEED ATTENTION:**

DISTRIBUTION

PARALLELISM

WAYS OF DIRECTING THE SEARCH STRATEGY

WAYS OF DIRECTING THE MAPPING OF  
EXECUTION ON HARDWARE MODULES (TO  
TOLERATE HARDWARE FAULTS)

## **FAILURE MODES OF PROLOG PROGRAMS.**

(FOR ANY CAUSE: PROGRAMMING ERROR OR FAILURE IN THE SUPPORT)

THE NORMAL OUTPUT OF A PROLOG PROCEDURE (PROGRAM) MAY BE:

- SUCCESS WITH INSTANTIATION OF 0 OR MORE VARIABLES
- FAILURE (NO ASSIGNMENT OF VALUES TO VARIABLES SATISFIES THE CALLING GOAL)

### **FAILURE MODES:**

- FAILURE TO PRODUCE A SOLUTION THAT EXISTS, WITH INFINITE SEARCH
- FAILURE TO PRODUCE A SOLUTION THAT EXISTS (I.E. IS LOGICALLY DEDUCTIBLE FROM THE DATABASE)
- PRODUCTION OF A WRONG SOLUTION (LOGICAL ERROR)
- FAILURE TO PRODUCE A SOLUTION, BECAUSE IT IS NOT DEDUCTIBLE FROM THE DATABASE (LOGICAL ERROR)
- NON-DETERMINISM OF SOLUTION FROM REDUNDANT DATABASE

INVOLUNTARY/UNPREDICTABLE REDUNDANCY  
MAY ARISE FROM SELF MODIFYING  
PROGRAMS AS WELL AS FROM COMPLEXITY,  
E.G. THROUGH MAINTAINANCE

## PROBLEMS WITH THE LOGICAL PARADIGM:

- CLOSED WORLD ASSUMPTION;
- LACK OF NEGATION

## PROBLEMS WITH THE SEARCH STRATEGY: VULNERABILITY TO

- UNBOUNDED RECURSION
- CIRCULAR REFERENCES

```
ancestor(X,Y) :- parent(X,Y).
ancestor(X,Y) :- ancestor(X,Z), ancestor(Z,Y).
parent(abraham, isaac).
parent(isaac, jacob).
```

```
/* The query
?- ancestor(A,B).
```

```
obtains the answers:
  A = abraham, B = isaac ;
  A = isaac, B = jacob ;
  A = abraham, B = jacob ;
Not enough heap space.
*/
```

SOLUTIONS FOR THIS TYPE OF FAILURE:  
LIMITING DEPTH OF SEARCH; PARALLEL  
SEARCH

OF COURSE, THIS TOY EXAMPLE WOULD  
WORK CORRECTLY IF WRITTEN AS:

```
ancestor(X,Y) :- parent(X,Y).  
ancestor(X,Y) :- parent(X,Z), ancestor(Z,Y).  
parent(abraham, isaac).  
parent(isaac, jacob).  
parent(jacob,a).
```

```
/* The query:  
?- ancestor(A,B).  
gets the answers:  
  A = abraham, B = isaac ;  
  A = isaac, B = jacob ;  
  A = abraham, B = jacob ;  
no  
*/
```

## LANGUAGE FEATURES FOR IMPROVING RELIABILITY OF PROGRAMS:

NORMAL FEATURES OF HIGH-LEVEL LANGUAGES:

MODULES WITH NAME SCOPE RULES

ANNOTATION OF INPUT AND OUTPUT ARGUMENTS

LIMITATIONS TO SELF-MODIFIABILITY THROUGH ASSERT AND RETRACT

EXCEPTION HANDLING FOR:

- INTERPRETER-GENERATED EXCEPTIONS;
- PROGRAMMER-DEFINED EXCEPTIONS

OTHER FEATURES FOR ROBUSTNESS AT RUN-TIME :

LIMITS ON DEPTH OF SEARCH (LOOP DETECTION, TIME-OUTS OR NUMBER OF STEPS?)

PARALLEL SEARCH: SIMPLIFIES REDUNDANT PROGRAMMING, RISKS NON-DETERMINISM

LIMITS ON EFFECTS OF SELF-MODIFICATION (ATOMICITY RULES)?



## TENTATIVE CONCLUSIONS:

- PROLOG HAS PROBLEMS THAT CAN BE CORRECTED BY NORMAL SOFTWARE ENGINEERING REMEDIES: STRUCTURE AND DISCIPLINE
- THE NON-PROCEDURAL STYLE HAS PROS AND CONS:

SHIFT OF COMPLEXITY FROM APPLICATION TO INTERPRETER: +

LESS CONTROL ON PROGRAM BEHAVIOUR: -

NON-PORTABILITY OF PROGRAMS BETWEEN INTERPRETERS: -

SIMPLICITY OF ADDING REDUNDANCY: +

DIFFICULTY IN DIRECTING USE OF REDUNDANCY: -

- THERE ARE PECULIAR FAILURE MODES:

DIFFERENCE BETWEEN "FAILED" PROCEDURE  
ACTIVATION (NO SOLUTION) AND WRONG  
SOLUTION

UNCONTROLLED REDUNDANCY